The relatively young area of forgetting is concerned with the removal of selective information, while preserving other knowledge. This might be useful or even necessary, for example, to simplify a knowledge base or to tend legal requests. In the last few years, there has been an ample amount of research in the field, in particular with respect to logic programs, spanning from theoretical considerations to more practical applications, starting at the conceptual proposal of forgetting, to suggestions of properties that should be satisfied, followed by characterizations of abstract classes of operators that satisfy these properties, and finally the definition of concrete forgetting procedures.

In this work we present novel Python implementations of all the forgetting procedures that have been proposed to date on logic programs. We provide them in a web interface, and hope to thereby give anybody who is interested a low-barrier overview of the landscape.

An Overview

While dynamics in knowledge representation are often examined with respect to the addition of new information, there is an increasing amount of research on how it may be removed. Forgetting, or variable elimination, seeks to make a knowledge base independent of elements of the underlying formal language, while keeping as many logical connections between the remaining elements as possible. Though it was initially proposed as a semantical notion over classical formulas [LR94], in the context of logic programming forgotten atoms are usually also required to be removed syntactically. More broadly, there have been several suggestions for properties that a plausible forgetting procedure should satisfy [EW08, ZZ09, WZZZ12, WWZ13, KA14, DW15], cf. [GKL16b] for an extensive overview.

Along the way, there have also been a number of suggestions for concrete syntactical transformations to forget atoms from a program [ZF06, EW08, KA14, BGKL19, GJKL21, ACF+22b, ACF+22a, Ber22] that satisfy some of the proposed properties, or at least satisfy them whenever possible [GKL16a]. The hope for such operators is to produce forgetting results that in some way resemble their origin, while avoiding a computational blow-up as much as possible. Given a semantic definition of a class of forgetting operators, e.g. [GKLW17, GJK+19], the construction of a program from the desired HT-models remains a baseline that can be employed to confirm the existence of a concrete forgetting operator [CF07].

The definitions of the syntactic forgetting procedures are rather involved, sometimes spanning over several pages of text. In order to make them more accessible, we compiled encodings as well as short explanations, of the operators (those that are defined over logic programs) and the construction of canonical programs into a web interface.1 We hope to thereby give anybody who is interested a good overview of the material.

A screenshot of the interface is shown in Figure 1. In the leftmost column, users can specify a logic program and the atoms to be forgotten, or select a predefined sample input. Subsequently, they select the

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1ForgettingWeb can be accessed here. Its source code is available here.
ForgettingWeb

Figure 1: Screenshot of ForgettingWeb.

forgetting operator of their choice from the dropdown in the middle column. These are implementations from operators proposed in the literature. For more details, the user presses the "Operator Explanation" button: this triggers a pop-up with an explanation of the operator and a reference to the paper in which it was proposed. After pressing the "Forget it!" button, the result program appears in the rightmost column.

An input program $P$ is specified by one rule per line, where the head atoms are separated by ‘;’, the head and the body are separated by ‘:-’, and the body literals are separated by commas. As negation signs both the traditional ‘not’ and the shorter ‘∼’ can be used. No dot at the end of a rule is required.

The atoms of the forgetting set $V$ are separated by commas. In the “Result” column on the right side, the new program has the same syntax as the input program.

The input of the counter-models construction $\text{aux}_{cm}$ are tuples with delimiters ‘<’ and ‘>’, e.g.:

$$<ab,ab><a,ab><b,ab><,ab>$$

$$<a,a><,>$$

Each character within an element of a tuple is assumed to be a single element of a set.

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