ON LAX EPIMORPHISMS AND THE ASSOCIATED FACTORIZATION

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Abstract. We study lax epimorphisms in 2-categories, with special attention to Cat and $\mathcal{V}$-$\text{Cat}$. We show that any 2-category with convenient colimits has an orthogonal $\text{LaxEpi}$-factorization system, and we give a concrete description of this factorization in Cat.

1. Introduction

A morphism $e : A \to B$ in a category $\mathcal{A}$ is an epimorphism if, for every object $C$, the map $\mathcal{A}(e, C) : \mathcal{A}(B, C) \to \mathcal{A}(A, C)$ is injective; looking at the hom-sets as discrete categories, this means that the functor $\mathcal{A}(e, C)$ is fully faithful. Lax epimorphisms (also called co-fully-faithful morphisms) are a 2-dimensional version of epimorphisms; in a 2-category they are precisely the 1-cells $e$ making $\mathcal{A}(e, C)$ fully faithful for all $C$.

One of the most known (orthogonal) factorization systems in the category of small categories and functors is the comprehensive factorization system of Street and Walters [20]. Another known factorization system consists of bijective-on-objects functors in the left-hand side and fully faithful functors in the right. Indeed in both cases we have an orthogonal factorization system in the 2-category $\text{Cat}$ in the sense of Definition 3.1. This means that with the usual notion in ordinary categories we have a 2-dimensional aspect of the diagonal fill-in property. Here we show that $\text{Cat}$ has also an orthogonal $(\mathcal{E}, \mathcal{M})$-factorization system where $\mathcal{E}$ is the class of all lax epimorphisms, and present a concrete description of it, making use of a characterization of the lax epimorphic functors given in [1] (Theorem 4.5).

Moreover, any 2-category has an orthogonal $(\text{LaxEpi}, \text{LaxStrongMono})$-factorization system provided that it has 2-colimits and is almost cowellpowered with respect to lax epimorphisms (Theorem 3.16). Here to be almost cowellpowered with respect to a class $\mathcal{E}$ of morphisms means that, for every morphism $f$, the category of all factorizations $d \cdot e$ of $f$ with $e \in \mathcal{E}$ has a weakly terminal set. A key property is the fact that lax epimorphisms are closed under 2-colimits (Theorem 3.10).

We dedicate the last section to the study of lax epimorphisms in the 2-category $\mathcal{V}$-$\text{Cat}$ for $\mathcal{V}$ a complete symmetric monoidal closed category. In this context, it is natural to consider a variation of the notion of lax epimorphism: We say that a $\mathcal{V}$-functor $J : \mathcal{A} \to \mathcal{B}$ is a $\mathcal{V}$-lax epimorphism if the $\mathcal{V}$-functor $\mathcal{V}$-$\text{Cat}[J, \mathcal{C}] : \mathcal{V}$-$\text{Cat}[\mathcal{B}, \mathcal{C}] \to \mathcal{V}$-$\text{Cat}[\mathcal{A}, \mathcal{C}]$ is $\mathcal{V}$-fully faithful for all small $\mathcal{V}$-categories $\mathcal{C}$. Assuming that $\mathcal{V}$ is also cocomplete, Theorem 5.6 gives several characterizations of the lax epimorphisms in the 2-category $\mathcal{V}$-$\text{Cat}$. In particular, we show that they are precisely the $\mathcal{V}$-lax epimorphisms, and also precisely those $\mathcal{V}$-functors for which there is an isomorphism...
Moreover, \( \mathcal{V} \)-lax epimorphisms are equivalently defined if above we replace all small \( \mathcal{V} \)-categories \( C \) by all possibly large \( \mathcal{V} \)-categories \( C \), or by just the category \( \mathcal{V} \). This last characterization as well as Theorem 5.9, which characterizes \( \mathcal{V} \)-lax epimorphisms as absolutely \( \mathcal{V} \)-(co)dense \( \mathcal{V} \)-functors, have been proved before for \( \mathcal{V} = \text{Set} \) in [1].

For the basic theory on 2-categories we refer to [15] and [16]. For a detailed account on 2-dimensional (co)limits, see [14]; here we use the notation \( \lim(W, F) \) for the limit of \( F : \mathcal{A} \to \mathcal{B} \) weighted (“indexed” in Kelly’s language) by \( W : \mathcal{A} \to \text{Cat} \). Concerning enriched categories, we refer to [13].

2. Lax epimorphisms in 2-categories

In this section we present some basic properties and examples on lax epimorphisms. We end up by showing that, under reasonable conditions, for 2-categories \( \mathcal{S} \) and \( \mathcal{B} \), every lax epimorphism of the 2-category \( 2\text{-Cat}[\mathcal{S}, \mathcal{B}] \) is pointwise. Pointwise lax epimorphisms will have a role in the main result of Section 3.

**Definition 2.1.** A lax epimorphism in a 2-category \( \mathcal{A} \) is a 1-cell \( f : A \to B \) for which all the hom-functors
\[
\delta(f, C) : \delta(B, C) \to \delta(A, C)
\]
(with \( C \in \mathcal{A} \)) are fully faithful.

**Remark 2.2** (Duality and Coduality). The notion of lax epimorphism is dual to the notion of fully faithful morphism (in a 2-category). That is the reason why lax epimorphisms are also called co-fully-faithful morphisms. Indeed, the notion fully faithful morphism in the 2-category of small categories \( \text{Cat} \) coincides with the notion of fully faithful functor, since a functor \( P : \mathcal{A} \to \mathcal{B} \) is fully faithful if and only if the functor \( \text{Cat}(\mathcal{C}, P) : \text{Cat}(\mathcal{C}, \mathcal{B}) \to \text{Cat}(\mathcal{C}, \mathcal{A}) \) is fully faithful for all categories \( \mathcal{C} \).

On the other hand, the notion of lax epimorphism is self-codual. Namely, a morphism \( p : A \to B \) is a lax epimorphism in \( \mathcal{A} \) if and only if the corresponding morphism in \( \mathcal{A}^{\text{co}} \) (the 2-category obtained after inverting the directions of the 2-cells in \( \mathcal{A} \)) is a lax epimorphism.

**Remark 2.3.** Lax epimorphisms are closed for isomorphism classes. That is to say, if \( f \cong g \) and \( g \) is a lax epimorphism, then so is \( f \). Moreover, we have that lax epimorphisms are closed under composition and are right-cancellable: for composable morphisms \( r \) and \( s \), if \( r \) and \( sr \) are lax epimorphisms, so is \( s \).

**Examples 2.4.**

1. In a locally discrete 2-category, lax epimorphisms are just epimorphisms, since fully faithful functors between discrete categories are injective functions on the objects. But, in general, the class of lax epimorphisms and the one of epimorphisms are different and no one contains the other (see [1]).

2. Coequifiers are lax epimorphisms. The property of being a lax epimorphism encompasses the two-dimensional aspect of the universal property of a coequifier (see [14] pag. 309]). But, as observed in [1], coequalizers in \( \text{Cat} \) are not necessarily lax epimorphisms.

3. Any equivalence is a lax epimorphism. Recall that a morphism \( g : A \to B \) is an equivalence if there is \( f : B \to A \) with \( gf \cong 1_B \) and \( fg \cong 1_A \). This is equivalent to the existence of an adjunction between \( f \) and \( g \) with both unit and counit being invertible, and it is also well known that it is equivalent to the existence of an adjunction \( (\varepsilon, \eta) : f \dashv g \) together with
both $f$ and $g$ fully faithful. Dually, $g : A \to B$ is an equivalence if and only if there is an
adjunction $(\varepsilon, \eta) : f \dashv g$ with both $f$ and $g$ being lax epimorphisms. Moreover, given an
adjunction $(\varepsilon, \eta) : f \dashv g : A \to B$ in a 2-category $\mathcal{A}$, the morphism $g$ is a lax epimorphism if
and only if $f$ is fully faithful, if and only if $\eta$ is invertible (see \cite{16} Lemma 2.1.1).

(4) In a locally thin 2-category (i.e., with the hom-categories being preordered sets), the lax
epimorphisms are the order-epimorphisms, i.e., morphisms $f$ for which $g \cdot f \leq h \cdot f$ implies
$g \leq h$; and coinserters are lax epimorphisms – this immediately follows from the definition
of coinserter (see, for instance, \cite{14}, pag. 307).

However, coinserters are not lax epimorphisms in general; we indicate a simple counter-
example in the 2-category $\text{Cat}$ of small categories.\footnote{This rectifies \cite{11} Example 2.1.1.}
Let $A$ be the discrete category with a
unique object $A$, $B$ the discrete category with two objects, $FA$ and $GA$, and $F,G : A \to B$
the functors defined according to the name of the objects of $B$. The coinserter of $F$ and
$G$ is an inclusion $P : B \to C$, where $C$ has the same objects as $B$ and a unique non trivial
morphism, $\alpha_A : FA \to GA$. More precisely, the coinserter is given by the pair $(P, \alpha)$. (For
a description of coinserters in $\text{Cat}$, see \cite{5}, Example 6.5.) But $P$ is not a lax epimorphism.
Indeed, let $J, K : C \to D$ be two functors, where the category $D$ consists of four objects and
six non trivial morphisms as in the diagram below, with $KA \cdot \gamma_{FA} = r \neq s = \gamma_{GA} \cdot JA:

\[
\begin{array}{ccc}
JFA & \xrightarrow{\gamma_{FA}} & KFA \\
\downarrow r & & \downarrow s \\
JGA & \xrightarrow{\gamma_{GA}} & KGA
\end{array}
\]

Then, we have a natural transformation $\gamma : JP \to KP$ which cannot be expressed as $\gamma = \overline{\gamma} \ast \text{id}_P$ for any $\overline{\gamma} : J \Rightarrow K$.

(5) In the 2-category $\text{Pos}$ of posets, monotone functions and pointwise order between maps,
lax epimorphisms coincide with epimorphisms, and also with coinserters of some pair of
morphism (see \cite{4} Lemma 3.6]).

(6) In $\text{Preord}$, lax epimorphisms need not to be epimorphisms: they are just the monotone
maps $f : A \to B$ such that every $b \in B$ is isomorphic to $f(a)$ for some $a$.

Moreover, coinserters are strictly contained in lax epimorphisms, they are precisely the
monotone bijections. Indeed, given $f,g : A \to B$, let $\bar{B}$ be the underlying set of $B$ with the
preorder given by the reflexive and transitive closure of $\leq_B \cup \leq^\prime$, where $\leq_B$ is the order in
$B$ and $\gamma \leq^\prime z$ whenever there is some $x \in A$ with $y \leq f(x)$ and $g(x) \leq z$; the coinserter
is the identity map from $B$ to $\bar{B}$. Conversely, if $h : B \to C$ is a monotone bijection, it is the
coinserter of the projections $\pi_1, \pi_2 : P \to B$, where $P$ is the comma object of $h$ along itself.

Observe that the functor $P : B \to C$ of Example \cite{4} is indeed a morphism of the full
2-subcategory $\text{Preord}$ of $\text{Cat}$; it is a lax epimorphism in $\text{Preord}$ but not in $\text{Cat}$.

(7) Let $\text{Grp}$ be the 2-category of groups, homomorphisms, and 2-cells from $f$ to $g$ in $\text{Grp}(A,B)$
given by those elements $\alpha$ of $B$ with $f(x) \circ \alpha = \alpha \circ g(x)$, for all $x \in A$ (where $\circ$ denotes
the group multiplication). The horizontal composition of $\alpha : f \to g$ with $\beta : h \to k : B \to C$ is
given by $\beta \ast \alpha = h(\alpha) \circ \beta = \beta \circ k(\alpha)$; and the unit on an arrow $f : A \to B$ is simply the neutral
element of $B$ (see \cite{6}).\footnote{This 2-category of all groupoids with just one object.}
The lax epimorphisms of \(\text{Grp}\) are precisely the regular epimorphisms, that is, surjective homomorphisms. Indeed, given a surjective homomorphism \(f : A \rightarrow B\), homomorphisms \(g, h : B \rightarrow C\) and an element \(\gamma \in C\), the equalities \(g(f(x)) \circ \gamma = \gamma \circ (h(f(x)))\) for all \(x \in A\) imply \(g(\gamma) \circ \gamma = \gamma \circ (h(\gamma))\) for all \(\gamma \in B\), showing that \(f\) is a lax epimorphism. Conversely, given a lax epimorphism \(f : A \rightarrow B\), consider its \((\text{RegEpi}, \text{Mono})\)-factorization in \(\text{Grp}\):

\[
\begin{array}{ccc}
A & \xrightarrow{q} & M \\
\downarrow^{m} & & \downarrow^{\mu} \\
B & & B
\end{array}
\]

Since \(q\) and \(qm\), so is \(m\), by Remark 2.3 We show that then \(m\) is an isomorphism. In \(\text{Grp}\), monomorphisms are regular (see [2]); let \(g, h : B \rightarrow C\) be a pair whose equalizer is the inclusion \(m : M \rightarrow B\), that is, \(M = \{y \in B \mid g(y) = h(y)\}\). Denoting the neutral element of \(C\) by \(e\), we have a 2-cell \(e : gm \rightarrow hm\). Since \(m\) is a lax epimorphism, there is a unique \(\alpha : g \rightarrow h\) with \(\alpha e = e\). But \(\alpha e = g(e) \circ \alpha = \alpha \circ h(e) = \alpha\); hence \(\alpha = e\), that is, \(g(y) \circ e = e \circ h(y)\) for all \(y \in B\). Thus, \(B = M\) and \(m\) is the identity morphism.

**Remark 2.5.** In [1], lax epimorphisms were characterized in the 2-category \(\text{Cat}\) of small categories, functors and natural transformations: given a functor \(F : A \rightarrow B\) and a morphism \(g : b \rightarrow c\) in \(B\), let \(g//F\) denote the category whose objects are triples \((h, a, k)\) such that the composition \(b \xrightarrow{h} Fa \xrightarrow{k} c\) is equal to \(g\), and whose morphisms \(f : (h, a, k) \rightarrow (h', a', k')\) are those \(f : a \rightarrow a'\) of \(A\) with \(Fa \cdot h = h'\) and \(k' \cdot Fa = k\). Then:

**Theorem 2.6.** [1] A functor \(F : A \rightarrow B\) is a lax epimorphism in \(\text{Cat}\) if and only if, for every morphism \(g\) of \(B\), the category \(g//F\) is connected.

**Remark 2.7.** Recall that a 2-functor \(G : \mathcal{A} \rightarrow \mathcal{B}\) is locally fully faithful if, for any \(A, B \in \mathcal{A}\), the functor \(G_{A,B} : \mathcal{A}(A, B) \rightarrow \mathcal{B}(G(A), G(B))\) is fully faithful.

It is natural to consider lax epimorphisms in the context of 2-adjunctions or biadjunctions. Let \((\varepsilon, \eta) : F \dashv G : \mathcal{A} \rightarrow \mathcal{B}\) be a 2-adjunction (respectively, biadjunction). In this case, we have that, for any \(A, B \in \mathcal{A}\),

\[
\begin{array}{ccc}
\mathcal{A}(A, B) & \xrightarrow{G_{A,B}} & \mathcal{B}(G(A), G(B)) \\
\downarrow^{\chi_{G(A,B)}} & & \downarrow^{X_{G(A,B)}} \\
\mathcal{A}(\varepsilon_A, B) & & \mathcal{A}(FG(A), B)
\end{array}
\]

(2.0.1)

commutes (respectively, commutes up to an invertible natural transformation), in which

\[
\chi_{G(A,B)} : \mathcal{B}(G(A), G(B)) \rightarrow \mathcal{A}(FG(A), B)
\]

\[
h \mapsto \varepsilon_B \circ F(h)
\]

is the invertible functor (respectively, equivalence) of the 2-adjunction (biadjunction).

In the situation above, since isomorphisms (respectively, equivalences) are fully faithful and fully faithful functors are left-cancellable (see Remark 2.3), we have that \(G_{A,B} : \mathcal{A}(A, B) \rightarrow \mathcal{B}(G(A), G(B))\) is fully faithful if, and only if, \(\mathcal{A}(\varepsilon_A, B)\) is fully faithful. Therefore, the 2-functor \(G : \mathcal{A} \rightarrow \mathcal{B}\) is locally fully faithful if and only if \(\varepsilon_C\) is a lax epimorphism for every \(C \in \mathcal{A}\).

**Remarks 2.8.** It is known that in a 2-category with cotensor products, fully faithful morphisms are those \(p : A \rightarrow B\) such that the comma object of \(p\) along itself is isomorphic to the cotensor
product $2 \otimes A$. Dually, assuming the existence of tensor products, a morphism $p : A \to B$ is a lax epimorphism if and only if

$$
\begin{array}{c}
A \xrightarrow{p} B \\
\downarrow \alpha \ast \text{id}_p \\
B \xrightarrow{v_0} \otimes B
\end{array}
$$

is an opcomma object, in which

$$
\begin{array}{c}
B \xrightarrow{v_1} \otimes B \\
\downarrow \alpha \\
B \xrightarrow{v_0}
\end{array}
$$

is the tensor product.

**Remarks 2.9** (Preservation and reflection of lax epimorphisms). Since, in the presence of tensor products, lax epimorphisms are characterized by opcomma objects as above, we conclude that:

**Lemma 2.10.** Let $F : \mathcal{B} \to \mathcal{A}$ be a 2-functor.

1. Assuming that $\mathcal{B}$ has tensor products, if $F$ preserves opcomma objects and tensor products, then $F$ preserves lax epimorphisms.
2. Assuming that $\mathcal{A}$ has tensor products, if $G$ creates opcomma objects and tensor products, then $G$ reflects lax epimorphisms.

Moreover, we also have that:

**Lemma 2.11.** Let $F \dashv G$ be a 2-adjunction.

1. The 2-functor $F : \mathcal{B} \to \mathcal{A}$ preserves lax epimorphisms.
2. If $G$ is essentially surjective, then $F$ reflects lax epimorphisms.

**Proof.** For any object $W$ of $\mathcal{A}$ and any morphism $p : A \to B$ of $\mathcal{B}$, we have that

$$
\begin{array}{c}
\mathcal{A}(F(B), W) \xrightarrow{\mathcal{A}(p, W)} \mathcal{A}(F(A), W) \\
\cong \mathcal{B}(B, G(W)) \xleftarrow{\mathcal{B}(p, G(W))} \mathcal{B}(A, G(W))
\end{array}
$$

(2.0.2)

commutes.

1. If $p : A \to B$ is a lax epimorphism in $\mathcal{B}$, for any $W \in \mathcal{A}$, we have that $\mathcal{B}(p, G(W))$ is fully faithful and, hence, by the commutativity of (2.0.2), $\mathcal{A}(F(p), W)$ is fully faithful.

2. Assuming that $G$ is essentially surjective, if $F(p) : F(A) \to F(B)$ is a lax epimorphism in $\mathcal{A}$, then, for any $Z \in \mathcal{B}$, there is $W \in \mathcal{A}$ such that $G(W) \cong Z$. Moreover, we have that $\mathcal{A}(F(p), W)$ is fully faithful and, hence, $\mathcal{B}(p, G(W))$ is fully faithful by the commutativity of (2.0.2). This implies that $\mathcal{B}(p, Z)$ is fully faithful for any $Z \in \mathcal{B}$. \qed
Definition 2.12. A 2-natural transformation $\lambda : F \to G : S \to B$ is:

1. a pointwise lax epimorphism if, for any $C \in S$, the morphism $\lambda_C : F(C) \to G(C)$ is a lax epimorphism in $B$;
2. a lax epimorphism if $\lambda$ is a lax epimorphism in the 2-category of 2-Cat$[S,B]$ of 2-functors, 2-natural transformations and modifications.

Proposition 2.13. Let $\lambda : F \to G : S \to B$ be a 2-natural transformation. If $\lambda$ is a pointwise lax epimorphism then it is a lax epimorphism in the 2-category 2-Cat$[S,B]$.

Proof. Let $\lambda : F \to G : A \to B$ be a 2-natural transformation with each $\lambda_A : FA \to GA$ a lax epimorphism in $B$. Let $\alpha, \beta : G \to H : A \to B$ be two 2-natural transformations, and let $\Theta : \alpha \ast \lambda \sim \beta \ast \lambda$ be a modification. In particular, we have 2-cells in $B$ indexed by $A \in A$:

$$
\begin{array}{cccc}
\lambda_A & GA & \alpha_A & \alpha_A \\
\downarrow \Theta_A & \downarrow & \downarrow & \downarrow \\
\lambda_A & HA & \beta_A & \beta_A \\
\end{array}
$$

This gives rise to unique 2-cells

$$
\begin{array}{cccc}
\lambda_A & GA & \alpha_A & \alpha_A \\
\downarrow \Theta_A & \downarrow & \downarrow & \downarrow \\
\lambda_A & HA & \beta_A & \beta_A \\
\end{array}
$$

with $\Phi_A \ast \lambda_A = \Theta_A$. The uniqueness of $\Phi = (\Phi_A)_{A \in A}$ is clear. It is straightforward to see that $\Phi$ is indeed a modification. \qed

However, not every lax epimorphism 2-natural transformation is a pointwise lax epimorphism. In fact, this is known to be true for epimorphisms and, as observed in (1) of Examples 2.4, lax epimorphisms in locally discrete 2-categories are the same as epimorphisms.

More precisely, consider the locally discrete 2-category $S$ generated by

$$
A \xrightarrow{h} B \xrightarrow{f} C
$$

with the equation $fh = gh$. The pair $(h,f)$ gives an epimorphism in 2-Cat$[2,S]$ but $h$ clearly is not an epimorphism in $S$. Since $S$ and 2-Cat$[2,S]$ are locally discrete, this proves that $(h,f)$ gives a 2-natural transformation which is a lax epimorphism but not a pointwise lax epimorphism.

Yet, it follows from Lemma 2.11 that the converse holds for many interesting cases. More precisely:

Theorem 2.14. Let $B$ be a 2-category with cotensor products. Then, a 2-natural transformation $\lambda : F \to G : S \to B$ is a lax epimorphism if and only if it is a pointwise lax epimorphism.

Proof. By Proposition 2.13, every pointwise lax epimorphism is an epimorphism. We prove the converse below.
Let 1 be the terminal category with only the object 0. For each \( s \in S \), we denote by \( 1 \to S \) the functor defined by \( s \). For each \( \overline{B} : 1 \to B \), we have the pointwise right Kan extension (see [8, Theorem I.4.2]) given by
\[
\text{Ran}_{\overline{B}}(a) = \lim \left( \mathfrak{S} \left( (a, s) \right), \overline{B} \right) \cong \mathfrak{S}(a, s) \cap \left( \overline{B} 0 \right).
\]
We conclude, then, that, for any \( s \in S \), we have the 2-adjunction
\[
2\text{-Cat}[S, B] \rightleftarrows \text{Ran}_{\overline{B}}.
\]
Therefore, by Lemma 2.11, assuming that \( \lambda : F \to G : S \to B \) is a lax epimorphism in \( 2\text{-Cat}[S, B] \), we have that, for every \( s \in S \),
\[
2\text{-Cat}[S, B](\lambda) = \lambda \ast \text{id}_{\overline{B}} = \lambda_{s}
\]
is a lax epimorphism in \( B \).

3. The orthogonal \textit{LaxEpi}-factorization system

Factorization systems in categories have largely shown their importance, taking the attention of many authors since the pioneering work exposed in [10]. (For a comprehensive account of the origins of the study of categorical factorization techniques see [21].) When the category has appropriate colimits, we get one of the most common orthogonal factorization systems, the \((\text{Epi}, \text{StrongMono})\) one. Since lax epimorphisms look an adequate 2-version of epimorphisms, it is natural to ask for a factorization system involving them. In this section, we will obtain an orthogonal \((\text{LaxEpi}, \text{LaxStrongMono})\)-factorization system in 2-categories. In the next section we give a description of this orthogonal factorization system in \( \text{Cat} \).

The notion of orthogonal factorization system in 2-categories generalizes the ordinary one by incorporating the two-dimensional aspect in the diagonal fill-in property. Here we use a strict version of the orthogonal factorization systems studied in [9] (see Remark 3.2):

**Definition 3.1.** In the 2-category \( \mathcal{A} \), let \( \mathcal{E} \) and \( \mathcal{M} \) be two classes of morphisms closed under composition and containing the isomorphisms. The pair \((\mathcal{E}, \mathcal{M})\) forms an orthogonal factorization system provided that:

(i) Every morphism \( f \) of \( \mathcal{A} \) factors as a composition \( f = me \) with \( e \in \mathcal{E} \) and \( m \in \mathcal{M} \).

(ii) For every \( A \xrightarrow{m} B \) in \( \mathcal{E} \) and \( C \xrightarrow{m} D \) in \( \mathcal{M} \), the square
\[
\begin{array}{ccc}
\mathcal{A}(B, C) & \xrightarrow{\mathcal{A}(B, m)} & \mathcal{A}(B, D) \\
\downarrow \mathcal{A}(e, C) & & \downarrow \mathcal{A}(e, D) \\
\mathcal{A}(A, C) & \xrightarrow{\mathcal{A}(A, m)} & \mathcal{A}(A, D)
\end{array}
\]
is a pullback in \( \text{Cat} \).

**Remark 3.2.** In [9], Dupont and Vitale studied orthogonal factorization systems in 2-categories in a non-strict sense. Thus, in (i) of Definition 3.1, the factorization holds up to equivalence, and in (ii), instead of a pullback, we have a bi-pullback.

**Remark 3.3.** The one-dimensional aspect of (ii) asserts, for each pair of morphisms \( f : A \to C \) and \( g : B \to D \) with \( mf = ge \), the existence of a unique \( t : B \to C \) with \( te = f \) and \( mt = g \). The two-dimensional aspect of (ii) means that, whenever, with the above equalities, we have \( t'e = f' \)
and \( mt' = g' \), and 2-cells \( \alpha : f \to f' \) and \( \beta : g \to g' \) such that \( m*\alpha = \beta*e \),

\[
\begin{array}{c}
A \xrightarrow{e} B \\
\downarrow f \quad \downarrow g \\
C \xrightarrow{\alpha} \downarrow m \quad \downarrow \beta \\
\downarrow t \quad \downarrow g' \\
D
\end{array}
\]

then there is a unique 2-cell \( \theta : t \to t' \) with \( \theta* e = \alpha \) and \( m*\theta = \beta \).

If \( E \) is made of lax epimorphisms, the two-dimensional aspect comes for free. Indeed, for \( \alpha : f = te \Rightarrow t'e = f' \), there is a unique \( \theta : t \Rightarrow t' \) with \( \theta*e = \alpha \); and, since \( \beta*e = m*\alpha = m*\theta*e \), we have \( \beta = m*\theta \).

**Definition 3.4.** A 1-cell \( m : C \to D \) is said to be a *lax strong monomorphism* if it has the diagonal fill-in property with respect to lax epimorphisms; that is, for every commutative square

\[
\begin{array}{c}
A \xrightarrow{e} B \\
\downarrow f \quad \downarrow g \\
C \xrightarrow{\alpha} \downarrow m \quad \downarrow \beta \\
\downarrow t \quad \downarrow g' \\
D
\end{array}
\]

with \( e \) a lax epimorphism, there is a unique \( t : B \to C \) such that \( te = f \) and \( mt = g \).

In other words, taking into account Remark 3.3, \( m : C \to D \) is a lax strong monomorphism if for every lax epimorphism \( e \), the morphisms \( e \) and \( m \) fulfil condition (ii) of Definition 3.1.

**Remark 3.5.** It is obvious that lax strong monomorphisms are closed under composition and left-cancelable; moreover, their intersection with lax epimorphisms are isomorphisms.

**Proposition 3.6.** In a 2-category:

(i) Every inserter is a lax strong monomorphism.

(ii) In the presence of coequifiers, every lax strong monomorphism is faithful, i.e., a morphism \( m : C \to D \) such that \( \Delta(X,m) \) is faithful for all \( X \).

**Proof.** (i) For the commutative square (3.0.2) above let \( e \) be a lax epimorphism and let the diagram

\[
\begin{array}{c}
D \xrightarrow{r} E \\
\downarrow a \quad \downarrow s \\
C \xleftarrow{m} \quad \downarrow a \\
D
\end{array}
\]

be an inserter. Since \( e \) is a lax epimorphism, there is a unique \( \beta : rg \Rightarrow sg \) with \( \alpha*f = \beta*e \). This implies the existence of a unique \( t : B \to C \) such that \( mt = g \) and \( \alpha*t = \beta \). Then we have \( \alpha*(te) = \beta*e = \alpha*f \) and \( m(te) = ge = mf \). Hence, by the universality of \((m,a)\), we conclude that \( te = f \). And \( t \) is unique: if \( mt = mt' \) and \( te = t'e \), then we have \( \alpha*t*e = \alpha*t'*e \), which implies \( \alpha*t = \alpha*t' \); this together with \( mt = mt' \) shows that \( t = t' \).

(ii) Given a lax strong monomorphism \( m : A \to B \) and two 2-cells \( \alpha, \beta : r \to s : X \to A \) with \( m*\alpha = m*\beta \), let \( e : A \to C \) be the coequifier of the 2-cells. Then \( m \) factors through \( e \). Since, by \( 2.4(2) \), \( e \) is a lax epimorphism, using the diagonal fill-in property, there is some \( t : C \to A \) with \( te = 1_A \). Then \( \alpha = \beta \). \qed
Examples 3.7. (1) In $\text{Pos}$ and $\text{Preord}$ the converse of \textbf{3.6(i)} also holds. In $\text{Pos}$ lax strong monomorphisms are just order-embeddings and order-embeddings coincide with inserters (\cite[Lemma 3.3]{4}).

Also in $\text{Preord}$ lax strong monomorphisms coincide with inserters. It is easily seen that lax strong monomorphisms are precisely the order-embeddings $m : X \to Y$ with $m[X]$ closed in $Y$ under isomorphic elements. Let $m : X \to Y$ be a lax strong monomorphism. Let $Z$ be obtained from $Y$ just replacing every element $y \in Y \setminus m[X]$ by two unrelated elements $(y, 1)$ and $(y, 2)$, and let the maps $f_1, f_2 : Y \to Z$ be equal on $m[X]$ and $f_i(y) = (y, i), i = 1, 2$, for the other cases. Endowing $Z$ with the least preorder which makes $f_1$ and $f_2$ monotone, we see that $m$ is the inserter of $f_1$ and $f_2$.

(2) But, in general, the converse of \textbf{3.6(i)} is false. Just consider an ordinary category (i.e. a locally discrete 2-category) where strong monomorphisms and regular monomorphisms do not coincide. For that, it suffices that regular monomorphisms are not closed under composition.

Remark 3.8. In contrast to \textbf{3.6} neither equifiers nor equalizers are, in general, lax strong monomorphisms. Consider the following equivalence of categories, where only the non trivial morphisms are indicated:

$$
\begin{array}{c}
\begin{array}{c}
A = a \\
E
\end{array}
\end{array}
\xrightarrow{f}
\begin{array}{c}
\begin{array}{c}
a \\
f \rightleftharpoons b
\end{array}
\end{array} = B
$$

The functor $E$ is a lax epimorphism (see Example \textbf{2.4(3)}), but not a lax strong monomorphism, since there is no $T : B \to A$ making the following two triangles

$$
\begin{array}{c}
A \\
E
\begin{array}{c}
\begin{array}{c}
\xrightarrow{E} B
\end{array}
\end{array}
\end{array}
\begin{array}{c}
\begin{array}{c}
A \\
E
\begin{array}{c}
\begin{array}{c}
\xrightarrow{T} B
\end{array}
\end{array}
\end{array}
\end{array}
$$

commutative. But $E$ is both an equifier and an equalizer. To see that it is an equalizer consider the pair $F, \text{id}_B : B \to B$, where $F$ takes all objects on $a$ and all morphisms on $1_a$. To see that it is an equafier consider the category

$$
(3.0.4)
$$

and 2-cells $\alpha, \beta : R \to S : B \to C$ given in the obvious way.

A key property in the sequel is the closedness of lax epimorphisms under colimits, in the sense of \textbf{3.9} below. The closedness of classes of morphisms under limits in ordinary categories was studied in \textbf{11}.

\footnote{A morphism $m : X \to Y$ in $\text{Pos}$ or $\text{Preord}$ is an order-embedding if $m$ is injective and $m(x) \leq m(y) \Leftrightarrow x \leq y$.}
Definition 3.9. Let $\mathcal{E}$ be a class of morphisms in a 2-category $\mathcal{B}$. We say that $\mathcal{E}$ is closed under (2-dimensional) colimits in $\mathcal{A}$ if, for every small 2-category $\mathcal{S}$, every weight $W : \mathcal{S}^{\text{op}} \to \text{Cat}$ and every 2-natural transformation $\lambda : D \to D' : \mathcal{S} \to \mathcal{B}$, the induced morphism

$$\text{colim}(W, \lambda) : \text{colim}(W, D) \to \text{colim}(W, D')$$

is a morphism in the class $\mathcal{E}$ whenever, for any $C \in \mathcal{S}$, $\lambda_C$ is a morphism in $\mathcal{E}$.

Theorem 3.10. Lax epimorphisms are closed under (2-dimensional) colimits.

Proof. In fact, if the 2-natural transformation $\lambda : D \to D' : \mathcal{S} \to \mathcal{B}$ is a pointwise lax epimorphism, then, for any $A \in \mathcal{B}$, the 2-natural transformation

$$\mathcal{B}(\lambda, A) : \mathcal{B}(D', A) \to \mathcal{B}(D, A),$$

pointwise defined by $\mathcal{B}(\lambda, A)_C = \mathcal{B}(\lambda_C, A)$, is pointwise fully faithful. Hence it is fully faithful in the 2-category $\text{Cat}[\mathcal{S}, \mathcal{B}]$ (dual of Proposition 2.13). Therefore, for any weight $W : \mathcal{S}^{\text{op}} \to \text{Cat}$ and $X \in \mathcal{B}$,

$$\mathcal{B}(\text{colim}(W, \lambda), X) \cong 2\text{-Cat}[\mathcal{S}, \mathcal{B}](W, \mathcal{B}(\lambda, A))$$

is fully faithful. This proves that $\text{colim}(W, \lambda)$ is a lax epimorphism in $\mathcal{B}$. \qed

Remark 3.11. As shown in [7], for any orthogonal $(\mathcal{E}, \mathcal{M})$-factorization system in an ordinary category, $\mathcal{E}$ and $\mathcal{M}$ are closed under, respectively, colimits and limits.

In $\text{Cat}$, lax epimorphisms are not closed under (2-dimensional) limits, fully faithful functors are not closed under (2-dimensional) colimits, and, moreover, equivalences are neither closed under limits nor colimits.

Indeed, consider the category $\triangledown 2$ with two objects and one isomorphism between them. Let $d^0$ and $d^1$ be the two possible inclusions $1 \to \triangledown 2$ of the terminal category in $\triangledown 2$. There is only one 2-natural transformation $\iota$ between the diagram

$$
\begin{array}{ccc}
1 & \xymatrix{d^0 \ar[r] & \triangledown 2 \ar[l]_d^1}
\end{array}
$$

and the terminal diagram $\begin{array}{ccc}1 & \xymatrix{1 \ar[l]}
\end{array}$.

Clearly, $\iota$ is a pointwise equivalence (and, hence, a pointwise lax epimorphism and fully faithful functor). However, the induced functor between the equalizers and the coequalizers are respectively

$$7 : \emptyset \to 1 \quad \text{and} \quad \iota : \Sigma \mathbb{Z} \to 1$$

in which $\Sigma \mathbb{Z}$ is just the group $(\mathbb{Z}, +, 0)$ seen as a category with only one object. The functor $7$ is not a lax epimorphism, while $\iota$ is not fully faithful. Hence, both are not equivalences.

Therefore, equivalences may not be the left or the right class of a (strict) orthogonal factorization system in a 2-category with reasonable (co)limits.

Remark 3.12. The closedness under colimits has several nice consequences, we indicate three of them, which are going to be useful in the proof of Theorem 3.16 below (cf. [11]).
(1) Let the two squares in the following picture be pushouts:

```
\begin{diagram}
  \node{A} \arrow{s,l}{id} \arrow{se}{f} \node{E} \arrow{n,l}{f} \\
  \node{A} \arrow{sw}{e_i} \node{E_i} \arrow{n,l}{e} \node{E} \arrow{s,l}{id} \\
\end{diagram}
```

Then, the dotted arrows form a 2-natural transformation between the corresponding origin diagrams, and the dashed arrow is the unique one induced by the universality of the inner square. From Theorem 3.10, if \( f \) is a lax epimorphism, so is \( f' \). In conclusion, lax epimorphisms are stable under pushouts.

(2) Analogously, we see that the cointersection \( e : A \to E \) of a family \( e_i : A \to E_i \) of lax epimorphisms is a lax epimorphism.

```
\begin{diagram}
  \node{A} \arrow{s,l}{id} \node{A} \arrow{se}{e_i} \node{E_i} \arrow{n,l}{e} \node{E} \\
\end{diagram}
```

(3) Moreover, the closedness under colimits ensures that, given a family of morphisms \( f_i : B \to C \) equalized by a lax epimorphism \( e \), i.e., \( f_i e = f_j e \) for all \( f_i \) and \( f_j \) of the family, the multiple coequalizer of all \( f_i \), if it exists, is also a lax epimorphism.

```
\begin{diagram}
  \node{E} \arrow{s,l}{f_i} \node{B} \arrow{n,l}{id} \node{B} \node{C} \arrow{s,l}{e} \\
\end{diagram}
```

**Remark 3.13.** Many of everyday categories are cowellpowered, that is, the family of epimorphisms with a same domain is essentially small. By contrast, in the “mother” of all 2-categories, Cat, the class of lax epimorphisms is not cowellpowered: For every cardinal \( n \), let \( A_n \) denote the category whose objects are \( a_i, i \in n \), and whose morphisms are \( f_{ij} : a_i \to a_j \) with \( f_{jk} f_{ij} = f_{ik} \) and \( f_{ii} = 1_{a_i} \) for \( i,j,k \in n \). Every inclusion functor \( E_n : A_0 \to A_n \), being an equivalence, is a lax epimorphism, but the family of all these \( E_n \) is a proper class. Moreover, the family \( E_n, n \in \text{Card}, \) fails to have a cointersection in Cat. However, Cat is almost cowellpowered in the sense of Definition 3.14 as shown in the next section.

**Definition 3.14.** Let \( \mathcal{E} \) be a class of 1-cells in a 2-category \( \mathcal{A} \). Given a morphism \( f : A \to B \), denote by \( \mathcal{E} \backslash f \) the category whose objects are factorizations \( A \xrightarrow{d} D \xrightarrow{p} B \) of \( f \) with \( d \in \mathcal{E} \), and whose morphisms \( u : (d,D,p) \to (e,E,m) \) are 1-cells \( u : D \to E \) with \( ud = e \) and \( md = p \). We say that \( \mathcal{A} \) is almost cowellpowered with respect to \( \mathcal{E} \), if \( \mathcal{E} \backslash f \) has a weakly terminal set for every morphism \( f \).

**Remark 3.15.** Clearly, a (2-)category with an orthogonal \( (\mathcal{E},\mathcal{M}) \)-factorization system is almost cowellpowered with respect to \( \mathcal{E} \): the \( (\mathcal{E},\mathcal{M}) \)-factorization of \( f : A \to B \) is indeed a terminal object of \( \mathcal{E} \backslash f \).
The closedness of lax epimorphisms under colimits allows to obtain the following:

**Theorem 3.16.** Let a 2-category \( \mathcal{A} \) have conical colimits and be almost cowellpowered with respect to lax epimorphisms. Then \( \mathcal{A} \) has an orthogonal \((\text{LaxEpi},\text{LaxStrongMono})\)-factorization system.

**Proof.** Let \( \mathcal{E} \) be the class of lax epimorphisms in \( \mathcal{A} \). Given a morphism \( f : A \rightarrow B \), let \( \{(e_i,E_i,m_i)\}_{i \in I} \) be a weakly terminal object of the category \( \mathcal{E}|f \); that is, for every factorization \( A \xrightarrow{d} D \xrightarrow{p} B \) of \( f \) with \( d \in \mathcal{E} \) there is some \( i \) and some morphism \( u : (d,D,p) \rightarrow (e_i,E_i,m_i) \). Take the cointersection \( e : A \rightarrow E \) of all \( e_i : A \rightarrow E_i \). By Remark 3.12(2), the morphism \( e \) belongs to \( \mathcal{E} \);

\[
\begin{array}{ccc}
A & \xrightarrow{e_i} & E_i \\
\downarrow e & & \downarrow m_i \\
E & & B
\end{array}
\]

moreover, the cointersection gives rise to a unique \( m : E \rightarrow B \) with \( me = f \). Thus, \( (e,E,m) \) is clearly a weakly terminal object of \( \mathcal{E}|f \).

Consider all \( s : E \rightarrow E \) forming a morphism \( s : (e,E,m) \rightarrow (e,E,m) \) in \( \mathcal{E}|f \). Let \( c : E \rightarrow C \) be the multiple coequalizer of the family of all these morphisms \( s : E \rightarrow E \). By Remark 3.12(3), \( c \) is a lax epimorphism.

Theorem 3.16. Let a 2-category \( \mathcal{A} \) have conical colimits and be almost cowellpowered with respect to lax epimorphisms. Then \( \mathcal{A} \) has an orthogonal \((\text{LaxEpi},\text{LaxStrongMono})\)-factorization system.

**Proof.** Let \( \mathcal{E} \) be the class of lax epimorphisms in \( \mathcal{A} \). Given a morphism \( f : A \rightarrow B \), let \( \{(e_i,E_i,m_i)\}_{i \in I} \) be a weakly terminal object of the category \( \mathcal{E}|f \); that is, for every factorization \( A \xrightarrow{d} D \xrightarrow{p} B \) of \( f \) with \( d \in \mathcal{E} \) there is some \( i \) and some morphism \( u : (d,D,p) \rightarrow (e_i,E_i,m_i) \). Take the cointersection \( e : A \rightarrow E \) of all \( e_i : A \rightarrow E_i \). By Remark 3.12(2), the morphism \( e \) belongs to \( \mathcal{E} \);

\[
\begin{array}{ccc}
A & \xrightarrow{e_i} & E_i \\
\downarrow e & & \downarrow m_i \\
E & & B
\end{array}
\]

moreover, the cointersection gives rise to a unique \( m : E \rightarrow B \) with \( me = f \). Thus, \( (e,E,m) \) is clearly a weakly terminal object of \( \mathcal{E}|f \).

Consider all \( s : E \rightarrow E \) forming a morphism \( s : (e,E,m) \rightarrow (e,E,m) \) in \( \mathcal{E}|f \). Let \( c : E \rightarrow C \) be the multiple coequalizer of the family of all these morphisms \( s : E \rightarrow E \). By Remark 3.12(3), \( c \) is a lax epimorphism. Since \( 1_E \) is one of those morphisms \( s \), and \( ms = m \) for all them, the universality of \( c \) gives a unique \( n : C \rightarrow B \) with \( nc = m \). It is easy to see that, \( c : (e,E,m) \rightarrow (ce,C,n) \) is also the coequalizer in \( \mathcal{E}|f \) of all the above morphisms \( s \).

We show that \( n : C \rightarrow B \) is a lax strong monomorphism. In the following diagram, let the outer square be commutative with \( q \in \mathcal{E} \); form the pushout \((\bar{q},\bar{r})\) of \( q \) along \( r \), and let \( w \) be the unique morphism with \( w\bar{q} = n \) and \( w\bar{r} = s \):

\[
(3.0.5)
\]

\[
\begin{array}{ccc}
P & \xrightarrow{q} & Q \\
\downarrow r & & \downarrow s \\
C & \xrightarrow{n} & B
\end{array}
\]

\[
\begin{array}{ccc}
P & \xrightarrow{q} & Q \\
\downarrow r & & \downarrow s \\
C & \xrightarrow{n} & B
\end{array}
\]

The closedness under colimits of lax epimorphisms ensures that \( \bar{q} \) is a lax epimorphism (Remark 3.12(1)), so \((\bar{q}ce,R,w) \in \mathcal{E}|f \). Since \((ce,C,n)\) is terminal, there is a unique \( u : R \rightarrow C \) forming a morphism in \( \mathcal{E}|f \); from \((\bar{q}ce,R,w)\) to \((ce,C,n)\), and it makes \( u\bar{q} : C \rightarrow C \) an endomorphism on \((ce,C,n)\), then \( u\bar{q} = 1_C \). The morphism \( t = u\bar{r} \) fulfils the equalities \( tq = r \) and \( nt = s \). Moreover \( t \) is unique; indeed, if \( t' \) is another morphism fulfilling the same equalities, let \( k \) be the coequalizer of \( t \) and \( t' \) and let \( p : K \rightarrow B \) be such that \( pk = n \). Again by Remark 3.12, \((kc,K,p)\) belongs to \( \mathcal{E}|f \).

Arguing as before for \( \bar{q} \), we conclude that \( k \) is a split monomorphism, then \( t = t' \).

Taking into account Remark 3.3, we conclude that we have indeed an orthogonal factorization system in the 2-category \( \mathcal{A} \).

**Remark 3.17.** In \cite{[9]}, an orthogonal factorization system \((\mathcal{E},\mathcal{M})\) which, as the \((\text{LaxEpi},\text{LaxStrongMono})\) one, has \( \mathcal{E} \) made of lax epimorphisms and \( \mathcal{M} \) made of faithful morphisms is said to be \((1,2)\)-proper.

**Examples 3.18.** Some of the well-known orthogonal factorization systems in ordinary categories are indeed of the \((\text{LaxEpi},\text{LaxStrongMono})\) type for convenient 2-cells. This is the case in the 2-categories \( \text{Pos} \) and \( \text{Grp} \). In \( \text{Pos} \) it is the usual orthogonal \((\text{Surjections},\text{Order-embeddings})\)-factorization
system. Analogously for the category Top of topological spaces and continuous maps, with 2-cells given by the pointwise specialization order, we obtain \((\text{Surjections, Embeddings})\). For the 2-category Grp, the \((\text{LaxEpi}, \text{LaxStrongMono})\) factorization is precisely the \((\text{RegEpi}, \text{Mono})\) one.

Recall that, for every category with an orthogonal factorization system \((\mathcal{E}, \mathcal{M})\), we have that \(\mathcal{M} = \mathcal{E}^\perp\), i.e., \(\mathcal{M}\) consists of all morphisms \(m\) fulfilling the diagonal fill-in property as in (3.0.2) of Definition 3.4. From the proof of Theorem 3.16, Remark 3.11 and Remark 3.15 it immediately follows that, more generally, we have the following:

**Theorem 3.19.** Let \(\mathcal{E}\) be a class of morphisms closed under composition and containing isomorphisms in a cocomplete category \(\mathcal{A}\). Then, \((\mathcal{E}, \mathcal{E}^\perp)\) forms an orthogonal factorization system if and only if \(\mathcal{A}\) is almost cowellpowered with respect to \(\mathcal{E}\) and \(\mathcal{E}\) is closed under colimits.

---

4. **The Lax Epi-factorization in \(\text{Cat}\)**

In this section we describe the orthogonal \((\text{LaxEpi}, \text{LaxStrongMono})\)-factorization system in the 2-category \(\text{Cat}\) of small categories, functors and natural transformations. All we do applies also to the bigger universe \(\text{CAT}\) of possibly large categories.

Let us recall, by the way, two well-known orthogonal factorization systems \((\mathcal{E}, \mathcal{M})\) in the category \(\text{Cat}\):

(a) \(\mathcal{E}\) consists of all functors bijective on objects and \(\mathcal{M}\) consists of all fully faithful functors.

(b) \(\mathcal{E}\) consists of all initial functors and \(\mathcal{M}\) consists of all discrete opfibrations; analogously, for final functors and discrete fibrations [20].

It is easy to see that in both cases, (a) and (b), the system \((\mathcal{E}, \mathcal{M})\) fulfils the two-dimensional aspect of the fill-in diagonal property, thus we have an orthogonal factorization system in the 2-category \(\text{Cat}\) as defined in 3.1.

We start by defining discrete splitting bifibrations. We will see that they are precisely the lax strong monomorphisms.

**Notation 4.1.** Recall from Remark 2.5 the definition of the category \(g//P\) for a functor \(P : \mathcal{A} \to \mathcal{B}\) and a morphism \(g : b \to c\) of \(\mathcal{B}\). For every decomposition of \(g\) of the form \(b \overset{t}{\to} Pe \overset{k}{\to} c\), we denote by \([(r, s)]\) the corresponding connected component. By composing a morphism \(t : d \to b\) with \(C = [(r, s)]\) we obtain \(C \cdot t = [(rt, s)]\), a connected component of \(tg//P\). Analogously, for the composition on the right hand side: for \(u : b \to c\), \(u \cdot C = [(h, uk)]\).

**Definition 4.2.** Let \(P : E \to B\) be a functor.

(a) A \(P\)-split consists of a factorization of an identity \(1_b\) of the form

\[
\begin{array}{ccc}
  b & \xrightarrow{h} & Pe \xrightarrow{k} & b \\
  & \searrow_{1_b} & \downarrow & \\
  & & (hk, 1_{Pc}) & \\
\end{array}
\]

with \([(1_{Pc}, hk)] = [(hk, 1_{Pc})]\).
(b) A P-split diagram is a rectangle

\[(4.0.1)\]

where \((h,k)\) and \((h',k')\) are P-splits such that \([h,gk] = [(h'g,k')]\) in \(g/P\). The wavy line in the middle of the rectangle indicates the existence of an appropriate P-zig-zag between \((h,gk)\) and \((h'g,k')\); that is, the existence of a finite number of morphisms \(h_i, k_i, f_i\) making the following diagram commutative:

\[(4.0.1)\]

(c) The functor \(P : E \to B\) is said to be a discrete splitting bifibration if, for every P-split diagram \(4.0.1\), there is a unique commutative rectangle in \(E\) of the form

\[(4.0.2)\]

whose image by \(P\) is the outer rectangle of \(4.0.1\). (That is, \(P_{x_0} = x\), for each letter \(x\) with \(x_0\) appearing in \(4.0.2\).)

Remark 4.3. If \(P\) is a discrete splitting bifibration, then it is clear that, for every P-split of \(1_b\),

\[b \to Pe \to b\]

there are unique morphisms \(h_0 : b_0 \to e\) and \(k_0 : e \to b_0\) such that \(Ph_0 = h\) and \(Pk_0 = k\).

Proposition 4.4. Every discrete splitting bifibration

(1) is faithful,
(2) is conservative, and
(3) reflects identities.
Proof. Let \( P : E \to B \) be a discrete splitting bifibration.

(1) For \( \begin{array}{c} f \\ g \end{array} \xrightarrow{\hspace{1cm}} \begin{array}{c} a \\ b \end{array} \) with \( Pf = Pg = x \), consider the following diagrams:

\[
\begin{array}{c}
Pa \quad Pa \quad Pa \\
\downarrow \quad \downarrow \quad \downarrow \\
Pf \quad Pf \quad Pf \\
\end{array}
\quad \quad \quad \quad \quad \begin{array}{c}
a \quad a \quad a \quad a \\
\downarrow \quad \downarrow \quad \downarrow \\
f \quad g \quad g \\
\end{array}
\]

The first one is a \( P \)-split rectangle and it is the image by \( P \) of the two last ones. Then \( f = g \).

(2) Let \( f : a \to b \) be such that \( Pf \) is an isomorphism in \( B \). Then we have a \( P \)-split diagram:

\[
\begin{array}{c}
Pb \quad Pb \quad Pb \\
(Pf)^{-1} \quad \downarrow Pf \quad (Pf)^{-1} \\
Pa \quad Pa \quad Pa \\
\end{array}
\]

Consequently, there is a unique \( t_0 : b \to a \) with \( Pt_0 = (Pf)^{-1} \). Since, by (1), \( P \) is faithful, \( t_0 \) is the inverse of \( f \).

(3) Let \( f : d \to e \) be such that \( Pf = 1_x \). By (2), \( f \) is an isomorphism. Concerning the diagrams

\[
\begin{array}{c}
x \quad Pe \quad x \\
\downarrow \quad \downarrow \quad \downarrow \\
Pf \quad Pf \quad Pf \\
\end{array}
\quad \quad \quad \quad \quad \begin{array}{c}
d \quad f \quad e \\
\downarrow \quad \downarrow \quad \downarrow \\
d \quad f \quad e \\
\end{array}
\]

the first one is a \( P \)-split rectangle which is the image by \( P \) of the two rectangles on the right hand side. Consequently, \( f = 1_d \).

\[\square\]

Theorem 4.5. For \( \mathcal{E} \) the class of lax epimorphisms and \( \mathcal{M} \) the class of discrete splitting bifibrations, \((\mathcal{E}, \mathcal{M})\) is an orthogonal factorization system in \( \text{Cat} \) (and also in \( \text{CAT} \)).

Proof. Along the proof we represent the categories by blackboard bold letters: \( \mathbb{A}, \mathbb{B}, \) etc.

(1) The factorization. Given a functor \( F : \mathbb{A} \to \mathbb{B} \), we define the category \( \mathbb{E} \) as follows:

- \( \text{ob} \mathbb{E} \): pairs \((b, B)\) where \( b \in \mathbb{B} \) and \( B \) is a connected component of the category \( 1_b // F \);
- \( \text{mor} \mathbb{E} \): all \((b, B) \xrightarrow{g} (c, C)\) with \( g : b \to c \) a morphism of \( \mathbb{B} \) and \( g \cdot B = C \cdot g \), see Notation 4.1.

The identities and composition are obvious.

Let

\[
\mathbb{E} \xrightarrow{P} \mathbb{B}
\]

be the obvious projection, and define

\[
\mathbb{A} \xrightarrow{E} \mathbb{E}
\]

by \( Ea = (Fa, Ca) \) where \( Ca \) is the connected component of \((1_{Fa}, 1_{Fa})\) in \( 1_{Fa} // F \), and \( E(a \xrightarrow{f} a') = (Fa, Ca) \xrightarrow{Ff} (Fa', Ca') \). \( \mathbb{E} \) is clearly well-defined and \( F = P \cdot E \).

(2) \( E \) is a lax epimorphism. Using the characterization of lax epimorphisms given in [1], see Remark 2.5 we need to show that, for every \((b, B) \xrightarrow{g} (d, D)\) in \( \mathbb{E} \), the category \( g // E \) is connected.
Given two factorizations \((u_i, E a_i, v_i)\), \(i = 1, 2\), of \(g\) in \(E\) as in the figure

\[
(4.0.3)
\]

\[
\begin{array}{c}
(b, B) \\
\downarrow g \\
(d, D)
\end{array}
\]

by the definition of morphisms in \(E\), we have the following equalities of connected components in \(g//F\) (see Notation 4.1): \(g \cdot B = v_1 u_1 \cdot B = v_1 \cdot C_{a_1} \cdot u_1 = v_1 \cdot [(u_1, 1_{F a_1})] = [(u_1, v_1)];\) and, analogously, \(g \cdot B = [(u_2, v_2)]\), showing that \([(u_1, v_1)] = [(u_2, v_2)]\) in \(g//F\); hence, \([(u_1, v_1)] = [(u_2, v_2)]\) also in \(g//E\).

(3) \(P\) is a discrete splitting bifibration.

(3a) First observe that, given two factorizations in \(E\) of a same morphism \(g\) of the form

\[
\begin{array}{c}
P(e, E) \\
\downarrow r \\
\downarrow s \\
\downarrow c
\end{array}
\]

if \((r, (e, E), s)\) and \((r', (e', E'), s')\) belong to the same connected component of \(g//P\), then also \(s \cdot E \cdot r = s' \cdot E' \cdot r'\) in \(g//F\). Indeed, a \(P\)-zig-zag connecting these two factorizations, as illustrated in the left hand side diagram below gives rise to an \(F\)-zig-zag connecting \(s \cdot E \cdot r\) to \(s' \cdot E' \cdot r'\) in \(g//F\), as indicated in the right hand side diagram, where \(E = [(h, a, k)], E' = [(h', a', k')]}\) and \(E_j = [(h_j, a_j, k_j)]\):

(3b) Let

\[
(4.0.4)
\]

\[
\begin{array}{c}
b \\
\downarrow g \\
u_2 \\
\downarrow c
\end{array}
\]

be a \(P\)-split diagram with \(D = [(h_1, a_1, k_1)]\) and \(E = [(h_2, a_2, k_2)]\). Let \(B\) and \(C\) be the connected components of \(1_b//F\) and \(1_c//F\) given, respectively, by

\[B = v_1 \cdot D \cdot u_1 = [(h_1 u_1, a_1, v_1 k_1)]\] and \[C = v_2 \cdot C \cdot u_2 = [(h_2 u_2, a_2, v_2 k_2)]\].
By (3a), since \( [(1_d, u_1v_1)] = [(u_1v_1, 1_d)] \) in \( u_1v_1//P \), we have that \( u_1v_1 \cdot D = D \cdot u_1v_1 \). Then \( u_1B = u_1v_1Du_1 = Du_1v_1u_1 = Du_1 \) and \( Bv_1 = v_1Du_1v_1 = v_1u_1v_1D = v_1D \), showing that \( (b, B) \xrightarrow{u_1} (d, D) \) and \( (d, D) \xrightarrow{v_1} (b, B) \) are morphisms in \( E \). And \( B \) is unique, because, if \( B' \) is a connected component of \( 1b//F \) such that \( u_1B' = Du_1 \) and \( v_1D = B'v_1 \), then \( B' = v_1u_1B' = v_1Du_1 = B \). Analogously for \( c \xrightarrow{u_2} P(e, E) \xrightarrow{v_2} c \).

It remains to show that \( g : (b, B) \rightarrow (c, C) \) is a morphism of \( E \). By (3a), the \( P \)-split diagram (4.0.4) gives rise to the following \( F \)-split diagram:

\[
\begin{array}{ccc}
b & \xrightarrow{u_1} & d \xrightarrow{h_1} Fa_1 \xrightarrow{k_1} d \xrightarrow{v_1} b \\
\downarrow{g} & & \downarrow{k} \\
c & \xrightarrow{u_2} & e \xrightarrow{h_2} Fa_2 \xrightarrow{k_2} e \xrightarrow{v_2} c
\end{array}
\]

That is,

\[ gv_1Du_1 = v_2Eu_2g \ \text{in} \ g//F. \]

Hence, by definition of \( B \) and \( C \),

\[ gB = Cg, \]

showing that \( g \) is a morphism in \( E \). Since \( (b, B) \) and \( (c, C) \) are unique, \( g \) is clearly unique too. In conclusion, we have a unique diagram of morphisms of \( E \) of the form

\[
\begin{array}{ccc}
(b, B) & \xrightarrow{u_1} & (d, D) \xrightarrow{v_1} (b, B) \\
\downarrow{g} & & \downarrow{g} \\
(c, C) & \xrightarrow{u_2} & (e, E) \xrightarrow{v_2} (c, C)
\end{array}
\]

whose image by \( P \) is the rectangle of (4.0.4).

(4) \( (E, M) \) fulfils the diagonal fill-in property. Let

\[
\begin{array}{ccc}
A & \xrightarrow{Q} & B \\
\downarrow{G} & & \downarrow{H} \\
C & \xrightarrow{M} & D
\end{array}
\]

be a commutative diagram where \( Q \) is a lax epimorphism and \( M \) is a discrete splitting bifibration.

(4a) We define \( T : B \rightarrow C \) as follows:

Given \( b \in B \), since \( Q \) is a lax epimorphism, the category \( b//Q \) is connected. Let \( B \) be the unique connected component of \( 1b//Q \), and let \( (h, a, k) \) be a representative of \( B \). It is a \( Q \)-split, since \( (1_{Qa}, hk) \) and \( (hk, 1_{Qa}) \) belong to the same connected component of \( hk//Q \). Hence,

\[ Hb \xrightarrow{Hh} MGa \xrightarrow{Hk} Hb \]

is an \( M \)-split in \( D \).

By Remark [4.3], since \( M \) is a discrete splitting bifibration, there are unique morphisms \( h_0 : b_0 \rightarrow Ga \) and \( k_0 : Ga \rightarrow b_0 \) with \( Mh_0 = Hh \) and \( Mk_0 = Hk \). We put

\[(4.0.5) \quad Tb = b_0.\]
We show that $b_0$ does not depend on the representative of $B$. Indeed, for another representative $(h', a', k')$, we have a $Q$-split diagram as on the left hand side of (4.0.6); by applying $H$, we get the $M$-split diagram on the right hand side:

\[
\begin{array}{ccc}
  b & \xrightarrow{h} & Qa \\
  \downarrow & & \downarrow \\
  b & \xrightarrow{k'} & Qa'
\end{array}
\quad
\begin{array}{ccc}
  Hb & \xrightarrow{Hh} & MGa \\
  \downarrow & & \downarrow \\
  Hb & \xrightarrow{Hk'} & MGa'
\end{array}
\]

By hypothesis, there is a unique diagram

\[
\begin{array}{ccc}
  b_1 & \xrightarrow{h_1} & Ga \\
  \downarrow & & \downarrow \\
  s & \rightarrow & s \\
  b_2 & \xrightarrow{h_2} & Ga'
\end{array}
\]

whose image by $M$ is the outside rectangle of the first diagram of (4.0.6). But $M$ reflects identities, by Proposition 4.4 Then $s$ is an identity and, taking into account the unicity of $b_0$ and $k_0$ above, it must be $b_1 = b_2 = b_0$ and $s = 1_{b_0}$.

Let $b \xrightarrow{g} c$ be a morphism in $B$. Since $Q$ is a lax epimorphism, there is some $Q$-split diagram of the form

\[
\begin{array}{ccc}
  b & \xrightarrow{h_1} & Qa_1 \\
  \downarrow & & \downarrow \\
  c & \xrightarrow{h_2} & Qa_2
\end{array}
\]

By applying $H$ to it, we obtain an $M$-split diagram:

\[
\begin{array}{ccc}
  Hb & \xrightarrow{Hh_1} & MGa_1 \\
  \downarrow & & \downarrow \\
  Hc & \xrightarrow{Hh_2} & MGa_2
\end{array}
\]

By hypothesis, there are unique morphisms

\[
\begin{array}{ccc}
  b_0 & \xrightarrow{i_1} & Ga_1 \\
  \downarrow & & \downarrow \\
  c_0 & \xrightarrow{i_2} & Ga_2
\end{array}
\]

making the diagram commutative and whose image by $M$ is the rectangle of (4.0.7). We put $Tg = g_0$.

Again, by the unicity, we know that $b_0$ and $c_0$ do not depend on the representative of $1_{b//Q}$ and $1_{c//Q}$. And the unicity of $g_0$ follows then from the faithfulness of $M$ (Proposition 4.4).

$T$ is clearly a functor, the preservation of identities and composition being obvious.
(4b) We show that $T$ satisfies the diagonal fill-in condition. Given $b \in \text{ob} \mathcal{B}$, $MTb = M_{b_0} = Hb$, by construction, and, analogously, $MTg = Hg$, for each $g \in \text{mor} \mathcal{B}$.

Given $f : a \to a'$ in $\mathcal{A}$, the $M$-split diagram

$$
\begin{array}{ccc}
HQa = MGa & \to & MGa \\
HQf = MGf & \downarrow & MGf \\
HQa' = MGa' & \to & MGa'
\end{array}
$$

ensures that $TQf = Gf$.

Finally, if $T' : \mathcal{B} \to \mathcal{C}$ is another functor such that $T'Q = G$ and $MT' = H$, we show that $T = T'$. Let $g : b \to d$ be a morphism of $\mathcal{B}$. The morphism $T(b \overset{g}{\to} d) = b_0 \overset{g_0}{\to} d_0$ is the unique one making part of a commutative rectangle as in (4.0.8) whose image by $M$ is the rectangle of the $M$-split diagram (4.0.7). But the image by $M$ of the rectangle $T'b \overset{T'h}{\to} Ga \overset{T'k}{\to} T'b$

$$
\begin{array}{ccc}
T'g & \downarrow & T'g \\
T'd \overset{T'h'}{\to} Ga' \overset{T'k'}{\to} T'd
\end{array}
$$

gives also the $M$-split diagram (4.0.7). Then $T'g = g_0 = Tg$. □

**Question 4.6.** Inserters in $\text{Cat}$ are discrete splitting bifibrations (by Proposition 3.6). We don’t know if the converse is true or not.

## 5. Lax epimorphisms in the enriched context

In this section we study lax epimorphisms in the enriched setting.

**Assumption 5.1.** Along the section $\mathcal{V} = (\mathcal{V}_0, \otimes, I)$ is a symmetric monoidal closed category with $\mathcal{V}_0$ complete.

We denote by $\mathcal{V}$-$\text{Cat}$ the 2-category of small $\mathcal{V}$-categories, $\mathcal{V}$-functors and $\mathcal{V}$-natural transformations.

Let $\mathcal{A}$ be a small $\mathcal{V}$-category, and $\mathcal{B}$ a (possibly large) $\mathcal{V}$-category. By abuse of language, we also denote by $\mathcal{V}$-$\text{Cat}(\mathcal{A}, \mathcal{B})$ the category of $\mathcal{V}$-functors from $\mathcal{A}$ to $\mathcal{B}$ and $\mathcal{V}$-natural transformations between them. Moreover, in this setting, the designation $\mathcal{V}$-$\text{Cat}[\mathcal{A}, \mathcal{B}]$ (or just $[\mathcal{A}, \mathcal{B}]$) represents the $\mathcal{V}$-category of $\mathcal{V}$-functors; thus, for any pair of $\mathcal{V}$-functors $F, G : \mathcal{A} \to \mathcal{B}$, the hom-object $\mathcal{V}$-$\text{Cat}[\mathcal{A}, \mathcal{B}](F, G)$ is given by the end

$$
\int_{A \in \mathcal{A}} B(FA, GA).
$$

Recall that a $\mathcal{V}$-functor $P : \mathcal{A} \to \mathcal{B}$ is $\mathcal{V}$-fully faithful (called just fully faithful in [13]) if the map $P_{A,A'} : \mathcal{A}(A, A') \to B(FA, PA')$ is an isomorphism in $\mathcal{V}_0$ for all $A, A' \in \mathcal{A}$.

Let $\mathcal{I}$ be the unit $\mathcal{V}$-category with one object $0$ and $\mathcal{I}(0,0) = I$. Given a $\mathcal{V}$-functor $P : \mathcal{A} \to \mathcal{B}$, the underlying functor of $P$ is denoted by $P_0 = \mathcal{V}$-$\text{Cat}(\mathcal{I}, P) : \mathcal{A}_0 \to \mathcal{B}_0$.

In general, we use the notations of [13]; concerning limits, we denote a weighted limit over a functor $F : \mathcal{D} \to \mathcal{V}$ with respect to a weight $W : \mathcal{D} \to \mathcal{V}$ by $\text{lim}(W, F)$ (called indexed limit and designated by $\{W, F\}$ in [13]).
Lemma 5.2. For a \( \mathcal{V} \)-functor \( P : A \to B \), consider the following conditions.

(a) \( P \) is \( \mathcal{V} \)-fully faithful.
(b) \( P_0 \) is fully faithful.
(c) The functor \( \text{Cat}(C, P_0) : \text{Cat}(C, A_0) \to \text{Cat}(C, B_0) \) is fully faithful for every (ordinary) category \( C \).
(d) The functor \( \mathcal{V}-\text{Cat}(C, P) : \mathcal{V}-\text{Cat}(C, A) \to \mathcal{V}-\text{Cat}(C, B) \) is fully faithful for every \( \mathcal{V} \)-category \( C \).
(e) The \( \mathcal{V} \)-functor \( \mathcal{V}-\text{Cat}[C, P] : \mathcal{V}-\text{Cat}[C, A] \to \mathcal{V}-\text{Cat}[C, B] \) is \( \mathcal{V} \)-fully faithful for every \( \mathcal{V} \)-category \( C \).

We have that

\[
\begin{array}{cccccc}
(a) & \iff & (c) & \iff & (d) & \iff & (b)
\end{array}
\]

The five conditions are equivalent whenever (i) \( P \) has a left or right \( \mathcal{V} \)-adjoint, or (ii) \( \mathcal{V} = \mathcal{V}_0(I, -) : \mathcal{V}_0 \to \text{Set} \) is conservative.

Proof. It is well-known that \( (a) \iff (b) \) in case we have (i) or (ii) [13, 1.3 and 1.11].
\( (b) \iff (c) \). It is just Remark 2.2.
\( (a) \implies (d) \). Given two \( \mathcal{V} \)-functors \( F, G : C \to A \), and a \( \mathcal{V} \)-natural transformation \( \beta : PF \to PG \), we want to show that there is a unique \( \mathcal{V} \)-natural transformation \( \alpha : F \to G \) with \( P\alpha = \beta \). Since \( P \) is \( \mathcal{V} \)-fully faithful, \( P_{A,B} \) is a \( \mathcal{V}_0 \)-isomorphism for all \( A, B \in A \). We just define \( \alpha : F \to G \) with each component \( \alpha_C \) given by

\[
\alpha_C = \left( I \xrightarrow{\beta_C} B(PFC, PG) \xrightarrow{(P_{F,G})^{-1}} A(FC, GC) \right).
\]

Clearly \( \beta_C = P\alpha_C \) for each \( C \), and \( \alpha \) is unique. From the \( \mathcal{V} \)-naturality of \( \beta \) and the fact that \( P \) is a \( \mathcal{V} \)-functor, it immediately follows that \( \alpha \) is \( \mathcal{V} \)-natural.
\( (d) \implies (b) \). It follows from the fact that \( P_0 = \mathcal{V}-\text{Cat}(I, P) \) by definition.
\( (a) \implies (c) \). Recall that there is a bijection

\[
A \ni A \mapsto \overline{A} \in \mathcal{V}-\text{Cat}[I, A]
\]

in which \( \overline{A} : I \to A \) is the only \( \mathcal{V} \)-functor from the unit \( \mathcal{V} \)-category \( I \) to \( A \) such that \( \overline{A}_0 = A \). Moreover, for any \( A, B \in A \), the hom-object \( A(A, B) \) is the end \( \int \mathcal{A}(\overline{A} - , B) \) which gives the hom-object \( \mathcal{V}-\text{Cat}[I, A](\overline{A}, B) \). We get that, for any \( \mathcal{V} \)-functor \( P : A \to B \), the morphism \( P_{A,B} \) is essentially \( \mathcal{V} \)-Cat[I, A](\overline{A}, B).

Therefore \( \mathcal{V}-\text{Cat}[I, P] \) is \( \mathcal{V} \)-fully faithful if and only if \( P \) is \( \mathcal{V} \)-fully faithful.
\( (a) \implies (e) \). Given a \( \mathcal{V} \)-category \( C \) and \( \mathcal{V} \)-functors \( F, G : C \to A \), we have that

\[
\mathcal{V}-\text{Cat}[C, P][F,G] : \mathcal{V}-\text{Cat}[C, A](F, G) \to \mathcal{V}-\text{Cat}[B, A](PF, PG)
\]

is, by definition, the morphism

\[
\int_{C \in C} P_{(FC, GC)} : \int_{C \in C} A(FC, GC) \to \int_{C \in C} B(PFC, PG).
\]

induced by the \( \mathcal{V} \)-natural transformation between the \( \mathcal{V} \)-functors \( A(F-, G-) \) and \( B(PF-, PG-) \) whose components are given by

\[
P_{F,A,GB} : A(FA, GB) \to B(PFA, PGB).
\]

Since \( P \) is \( \mathcal{V} \)-fully faithful, we have that \( 5.0.2 \) is invertible and, hence, \( 5.0.1 \) is invertible. \( \Box \)
Recall that the counit of an adjunction \( F_0 \dashv G_0 \) between ordinary categories is invertible if and only if there is any natural isomorphism between \( F_0 G_0 \) and the identity \[ \text{[12] Lemma 1.3}\] From Lemma \[5.2\] we obtain:

**Lemma 5.3.** Given a \( \mathcal{V} \)-adjunction \((\epsilon, \eta) : F \dashv G : A \to B\), the \( \mathcal{V} \)-functor \( G \) is \( \mathcal{V} \)-fully faithful if and only if there is any (ordinary) natural isomorphism
\[
F_0 G_0 \to \text{id}_{A_0}.
\]

**Proof.** By Lemma \[5.2\] \( G \) is \( \mathcal{V} \)-fully faithful if and only if \( G_0 \) is fully faithful. But \( G_0 \) is fully faithful in \( \text{Cat} \) if and only if the counit \( \epsilon_0 \) is invertible if and only if there is any (ordinary) natural isomorphism \( F_0 G_0 \to \text{id}_{A_0} \). \( \square \)

On one hand, following Definition \[2.1\] a \( \mathcal{V} \)-functor \( P : A \to B \) between small \( \mathcal{V} \)-categories is said a lax epimorphism in the 2-category \( \mathcal{V} \text{-Cat} \) if the (ordinary) functor
\[
\mathcal{V} \text{-Cat}(P, C) : \mathcal{V} \text{-Cat}(B, C) \to \mathcal{V} \text{-Cat}(A, C)
\]
is fully faithful, for all \( \mathcal{V} \)-categories \( C \). On the other hand, the notion of \( \mathcal{V} \)-fully faithful functor and Lemma \[5.2\] inspire the following definition.

**Definition 5.4.** A \( \mathcal{V} \)-functor \( J : A \to B \) (between small \( \mathcal{V} \)-categories) is a \( \mathcal{V} \)-lax epimorphism if, for any \( C \) in \( \mathcal{V} \text{-Cat} \), the \( \mathcal{V} \)-functor
\[
\mathcal{V} \text{-Cat}[J, C] : \mathcal{V} \text{-Cat}[B, C] \to \mathcal{V} \text{-Cat}[A, C]
\]
is \( \mathcal{V} \)-fully faithful.

**Assumption 5.5.** Until now, we are assuming that \( \mathcal{V}_0 \), and then also the \( \mathcal{V} \)-category \( \mathcal{V} \), is complete (Assumption \[5.1\]). From now on, we assume furthermore that \( \mathcal{V}_0 \) is also cocomplete.

**Theorem 5.6.** Given a \( \mathcal{V} \)-functor \( J : A \to B \) between small \( \mathcal{V} \)-categories \( A \) and \( B \), the following conditions are equivalent.

(a) \( J \) is a \( \mathcal{V} \)-lax epimorphism.
(b) \( J \) is a lax epimorphism in the 2-category \( \mathcal{V} \text{-Cat} \).
(c) The functor \( \mathcal{V} \text{-Cat}(J, V) : \mathcal{V} \text{-Cat}(B, V) \to \mathcal{V} \text{-Cat}(A, V) \) is fully faithful.
(d) The \( \mathcal{V} \)-functor \( \mathcal{V} \text{-Cat}[J, V] : \mathcal{V} \text{-Cat}[B, V] \to \mathcal{V} \text{-Cat}[A, V] \) is \( \mathcal{V} \)-fully faithful.
(e) There is a \( \mathcal{V} \)-natural isomorphism \( \text{Lan}_J B(B, J) \cong B(B, -) \) (\( \mathcal{V} \)-natural in \( B \in \mathcal{B}^{\text{op}} \)).
(f) The \( \mathcal{V} \)-functor \( \mathcal{V} \text{-Cat}[J, C] : \mathcal{V} \text{-Cat}[B, C] \to \mathcal{V} \text{-Cat}[A, C] \) is \( \mathcal{V} \)-fully faithful for every (possibly large) \( \mathcal{V} \)-category \( C \).

**Proof.** (a) \( \Rightarrow \) (b). It follows from the implication (a) \( \Rightarrow \) (b) of Lemma \[5.2\] Namely, given a (small) \( \mathcal{V} \)-category \( C \), since \( \mathcal{V} \text{-Cat}[J, C] \) is \( \mathcal{V} \)-fully faithful, we get that \( \mathcal{V} \text{-Cat}[J, C]_0 = \mathcal{V} \text{-Cat}(J, C) \) is fully faithful.

(b) \( \Rightarrow \) (c). Given any \( \mathcal{V} \)-functors \( F, G : B \to \mathcal{V} \), we denote by \( P : C \to \mathcal{V} \) the full inclusion of the (small) sub-\( \mathcal{V} \)-category of \( \mathcal{V} \) whose objects are in the image of \( F \) or in the image of \( G \).

---

4See \[12\] Lemma 1.3] or \[17\] for further results on non-canonical isomorphisms.

5Consider the diagram \[2.1.1\] for the case of adjunction between ordinary categories.
It should be noted that \( \mathcal{V}\text{-Cat}(J, \mathcal{C}) \) is a bijection by hypothesis, and \( \mathcal{V}\text{-Cat}(A, P), \mathcal{V}\text{-Cat}(B, P) \) are bijections since \( P \) is \( \mathcal{V} \)-fully faithful. Therefore, since the diagram

\[
\begin{array}{ccc}
\mathcal{V}\text{-Cat}(B, C)(F, G) & \xrightarrow{\mathcal{V}\text{-Cat}(J, C)_F} & \mathcal{V}\text{-Cat}(A, C)(F \cdot J, G \cdot J) \\
\mathcal{V}\text{-Cat}(B, P)_F & \downarrow & \mathcal{V}\text{-Cat}(A, P)_F \\\n\mathcal{V}\text{-Cat}(B, V)(F, G) & \xrightarrow{\mathcal{V}\text{-Cat}(J, V)_F} & \mathcal{V}\text{-Cat}(A, V)(F \cdot J, G \cdot J)
\end{array}
\]

commutes, we conclude that \( \mathcal{V}\text{-Cat}(J, V)_F \) is also a bijection. This proves that \( \mathcal{V}\text{-Cat}(J, V) \) is fully faithful.

\( \ominus \Rightarrow (\Box) \). Since \( \mathcal{V} \) is complete, we have that \( \mathcal{V}\text{-Cat}(J, \mathcal{V}) \) has a right \( \mathcal{V} \)-adjoint given by the (pointwise) Kan extensions \( \text{Ran}_J \). Therefore, assuming that \( \mathcal{V}\text{-Cat}(J, \mathcal{V}) \) is fully faithful, we conclude that \( \mathcal{V}\text{-Cat}(J, \mathcal{V}) \) is \( \mathcal{V} \)-fully faithful by Lemma 5.2.

\( (\Delta) \Rightarrow (\Box) \). Since \( \mathcal{V} \) is cocomplete, we have that \( \text{Lan}_J \dashv \mathcal{V}\text{-Cat}(J, \mathcal{V}) \). Therefore, assuming that \( \mathcal{V}\text{-Cat}(J, \mathcal{V}) \) is \( \mathcal{V} \)-fully faithful, we have the \( \mathcal{V} \)-natural isomorphism \( \epsilon : \text{Lan}_J(\cdot \cdot) \cong \text{id}_{\mathcal{V}\text{-Cat}(B, \mathcal{V})} \) given by the counit.

Denoting by \( \mathcal{V}_{Bop} \) the Enriched Yoneda Embedding (see, for instance, [13 2.4]), we have that \( \epsilon^{-1} \circ \text{id}_{\mathcal{V}_{Bop}} \) gives an isomorphism \( \text{Lan}_J B(B, J) \cong B(B, -) \) (\( \mathcal{V} \)-natural in \( B \in B^{op} \)).

\( (\Box) \Rightarrow (\Delta) \). Let \( \mathcal{C} \) be any (possibly large) \( \mathcal{V} \)-category. We consider the \( \mathcal{V} \)-functor \( \mathcal{V}\text{-Cat}[J, \mathcal{C}] \) and its factorization

\[
\begin{array}{ccc}
\mathcal{V}\text{-Cat}[B, \mathcal{C}] & \xrightarrow{\mathcal{V}\text{-Cat}[J, \mathcal{C}]_{\text{Im}}} & \text{Im}(\mathcal{V}\text{-Cat}[J, \mathcal{C}]) \\
& \downarrow & \text{Im}(\mathcal{V}\text{-Cat}[J, \mathcal{C}]) \\
& \downarrow & \mathcal{V}\text{-Cat}[A, \mathcal{C}] \\
& \downarrow & \mathcal{V}\text{-Cat}[J, \mathcal{C}]
\end{array}
\]

into a bijective on objects \( \mathcal{V} \)-functor \( \mathcal{V}\text{-Cat}[J, \mathcal{C}]_{\text{Im}} \) and the \( \mathcal{V} \)-full inclusion \( \text{Im}(\mathcal{V}\text{-Cat}[J, \mathcal{C}]) \rightarrow \mathcal{V}\text{-Cat}[A, \mathcal{C}] \) of the sub-\( \mathcal{V} \)-category \( \text{Im}(\mathcal{V}\text{-Cat}[J, \mathcal{C}]) \) whose objects are in the image of \( \mathcal{V}\text{-Cat}[J, \mathcal{C}] \). We prove below that \( \mathcal{V}\text{-Cat}[J, \mathcal{C}] \) is \( \mathcal{V} \)-fully faithful by proving that \( \mathcal{V}\text{-Cat}[J, \mathcal{C}]_{\text{Im}} \) is \( \mathcal{V} \)-fully faithful.

Given any \( \mathcal{V} \)-functor \( G : A \rightarrow C \) in \( \text{Im}(\mathcal{V}\text{-Cat}[J, \mathcal{C}]) \), we have that \( G = FJ \) for some \( F : B \rightarrow C \). Since \( \text{Lan}_J B(B, J) \equiv B(B, -) \), we conclude that \( \text{lim} \{ \text{Lan}_J B(B, J), F \} \) exists and, moreover, we have the isomorphisms

\[
\text{lim} \{ \text{Lan}_J B(B, J), F \} \cong \text{lim} (B(B, -), F) \cong F(B)
\]

by the (strong) Enriched Yoneda Lemma (see [13 Sections 2.4 and 4.1]).

Since \( \text{lim} \{ \text{Lan}_J B(B, J), F \} \) exists, it follows as a consequence of the universal property of left Kan extensions that \( \text{lim} (B(B, J), F \cdot J) \) exists and is isomorphic to \( \text{lim} \{ \text{Lan}_J B(B, J), F \} \) (see [13 Proposition 4.57]). Therefore, by (5.0.5) and by the formula for pointwise right Kan extensions (see [8 Theorem I.4.2] or, for instance, [13 Theorem 4.6]), we conclude that \( \text{Ran}_J (F \cdot J) \) exists and we have the isomorphism

\[
\text{Ran}_J (F \cdot J) B \cong \text{lim} (B(B, J), F \cdot J) \cong \text{lim} \{ \text{Lan}_J B(B, J), F \} \cong \text{lim} (B(B, -), F) \cong F(B)
\]
This proves that \( \text{Ran}_J(F \cdot J) \) exists for any \( G = F \circ J \) in \( \text{Im}(\mathcal{V}\text{-Cat}[J,C]) \), we conclude that \( \mathcal{V}\text{-Cat}[J,C]_{\text{Im}} \) has a right \( \mathcal{V} \)-adjoint, which we may denote by \( \text{Ran}_J \) by abuse of language. Finally, by the natural isomorphism \((5.0.6)\) and Lemma \(5.3\) we conclude that \( \mathcal{V}\text{-Cat}[J,C]_{\text{Im}} \) is \( \mathcal{V} \)-fully faithful.

\[ \text{(i)} \Rightarrow \text{(ii)}. \text{Trivial.} \]

\begin{remark}
For \( \mathcal{V} = \text{Set} \), the equivalence \((b) \Leftrightarrow (c)\) of Theorem \(5.6\) was given in \([1]\).
\end{remark}

\begin{remark} \textbf{(Duality).} \label{rem:duality}
A morphism \( J : A \to B \) is a lax epimorphism in \( \mathcal{V}\text{-Cat} \) if and only if \( J^{\text{op}} : A^{\text{op}} \to B^{\text{op}} \) is a lax epimorphism in \( \mathcal{V}\text{-Cat} \) as well.

Indeed, since the 2-functor \( \text{op} : \mathcal{V}\text{-Cat} \to \mathcal{V}\text{-Cat}^{\text{co}} \) is invertible, it takes lax epimorphisms to lax epimorphisms. Thus, \( J \) is a lax epimorphism in \( \mathcal{V}\text{-Cat} \) if, and only if, \( \text{op}(J) \) is a lax epimorphism in \( \mathcal{V}\text{-Cat}^{\text{co}} \) which, by Remark \(2.2\) holds if and only if \( J^{\text{op}} \) is a lax epimorphism in \( \mathcal{V}\text{-Cat} \).

Therefore, assuming that \( \mathcal{V}_0 \) is complete and cocomplete, \( J \) is a \( \mathcal{V} \)-lax epimorphism \( \iff J^{\text{op}} \) is a \( \mathcal{V} \)-lax epimorphism by Theorem \(5.6\).

Recall that a \( \mathcal{V} \)-functor \( J : A \to B \) between small \( \mathcal{V} \)-categories is \( \mathcal{V} \)-\textit{dense} if and only if its density comonad \( \text{Lan}_J \) is isomorphic to the identity on \( A \) (see \([13, \text{Theorem 5.1}]\)). Dually, \( J \) is \( \mathcal{V} \)-\textit{codense} if and only if the right Kan extension \( \text{Ran}_J \) is the identity. (Several concrete examples of (\( \mathcal{V} \)-)codensity monads are given in \([3]\).)

We say that \( J \) is \textit{absolutely} \( \mathcal{V} \)-\textit{dense} if it is \( \mathcal{V} \)-dense and \( \text{Lan}_J \) is preserved by any \( \mathcal{V} \)-functor \( F : B \to \mathcal{V} \). Dually, we define absolutely \( \mathcal{V} \)-\textit{codense} \( \mathcal{V} \)-\textit{functor}.

The following characterization of lax epimorphisms as absolutely dense functors was given in \([1]\) for \( \mathcal{V} = \text{Set} \):

\begin{theorem}
\label{thm:abs-dense}
Given a \( \mathcal{V} \)-functor \( J : A \to B \) between small \( \mathcal{V} \)-categories \( A \) and \( B \), the following conditions are equivalent.

\begin{enumerate}[(a)]
\item \( J \) is a \( \mathcal{V} \)-lax epimorphism.
\item \( J \) is absolutely \( \mathcal{V} \)-dense.
\item \( J \) is absolutely \( \mathcal{V} \)-codense.
\end{enumerate}
\end{theorem}

\begin{proof}
\text{(a)} \Rightarrow \text{(b)}. Assume that \( J \) is a \( \mathcal{V} \)-lax epimorphism. By \text{(c)} of Theorem \(5.6\) we have that \( B(B,-) \cong \text{Lan}_J B(B,-) \). Hence, since \( \text{lim}(B(B,-),\text{id}_B) \cong B \) exists by the (strong) Enriched Yoneda Lemma, we have that \( \text{lim}(\text{Lan}_J B(B,-),\text{id}_B) \) exists and is isomorphic to \( \text{lim}(B(B,-),\text{id}_B) \cong B \) (in which isomorphisms are always \( \mathcal{V} \)-natural in \( B \)).

Moreover, from the existence of \( \text{lim}(\text{Lan}_J B(B,-),\text{id}_B) \), we get that \( \text{lim}(B(B,-),J) \) exists and is isomorphic to \( \text{lim}(\text{Lan}_J B(B,-),\text{id}_B) \) (see \([13, \text{Proposition 4.57}]\)).

Finally, then, from the formula for pointwise right Kan extensions and the above, we get the \( \mathcal{V} \)-natural isomorphisms (in \( B \in B \))

\[ B \cong \text{lim}(B(B,-),\text{id}_B) \cong \text{lim}(\text{Lan}_J B(B,-),\text{id}_B) \cong \text{lim}(B(B,-),J) \cong \text{Ran}_J(B). \]

This proves that \( \text{Ran}_J \) is the identity on \( B \). That is to say, \( J \) is \( \mathcal{V} \)-codense.

Moreover, assuming that \( J \) is a \( \mathcal{V} \)-lax epimorphism, by Remark \(5.8\) \( J^{\text{op}} \) is a \( \mathcal{V} \)-lax epimorphism and, hence, by the proved above, \( J^{\text{op}} \) is \( \mathcal{V} \)-codense. Therefore \( J \) is \( \mathcal{V} \)-dense.
\end{proof}
By (d) of Theorem 5.6 we have that $\mathcal{V}$-$\text{Cat}[J,\mathcal{V}]$ is $\mathcal{V}$-fully faithful. Since $\mathcal{V}$ is cocomplete, we get that $\text{Lan}_J$ exists and there is an isomorphism $\text{Lan}_J(F \cdot J) \cong F$, $\mathcal{V}$-natural in $F \in \mathcal{V}$-$\text{Cat}[B,\mathcal{V}]$, given by the counit of $\text{Lan}_J \dashv \mathcal{V}$-$\text{Cat}[J,\mathcal{V}]$. This shows that $\text{Lan}_J$ is preserved by any $\mathcal{V}$-functor $F : B \to \mathcal{V}$.

(b) ⇒ (a). Assume that $J$ is absolutely $\mathcal{V}$-dense. We conclude that there is a natural isomorphism $\text{Lan}_J(F \cdot J) \cong F$. Therefore, by Lemma 5.2 we conclude that $\mathcal{V}$-$\text{Cat}[J,\mathcal{V}]$ is $\mathcal{V}$-fully faithful. By Theorem 5.6 this proves that $J$ is a $\mathcal{V}$-lax epimorphism.

(a) ⇔ (c). By Remark 5.8 and by the proved above, we conclude that $J$ is a $\mathcal{V}$-lax epimorphism $\iff J^{\text{op}}$ is absolutely $\mathcal{V}$-dense $\iff J$ is absolutely $\mathcal{V}$-codense.

□

Remark 5.10. Of course, density and codensity are not enough for a functor to be a lax epimorphism: for $1$ the terminal object in $\text{Cat}$, the functor $J : 1 \sqcup 1 \to 1$ is dense and codense, but not a lax epimorphism. Moreover, $\text{Ran}_J$ (respectively, $\text{Lan}_J$) is preserved by $F : 1 \to \text{Set}$ if and only if the image of $F$ is a preterminal object, i.e. the terminal set $\emptyset$; see [18, Remark 4.14] and [17, Remark 4.5].

References

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