



Contents lists available at ScienceDirect

European Journal of Combinatorics

journal homepage: www.elsevier.com/locate/ejc

The acquaintance time of (percolated) random geometric graphs

Tobias Müller^a, Paweł Prałat^b^a Mathematical Institute, Utrecht University, Utrecht, The Netherlands^b Department of Mathematics, Ryerson University, Toronto, ON, Canada

ARTICLE INFO

Article history:

Available online 7 March 2015

ABSTRACT

In this paper, we study the acquaintance time $\mathcal{AC}(G)$ defined for a connected graph G . We focus on $\mathcal{G}(n, r, p)$, a random subgraph of a random geometric graph in which n vertices are chosen uniformly at random and independently from $[0, 1]^2$, and two vertices are adjacent with probability p if the Euclidean distance between them is at most r . We present asymptotic results for the acquaintance time of $\mathcal{G}(n, r, p)$ for a wide range of $p = p(n)$ and $r = r(n)$. In particular, we show that with high probability $\mathcal{AC}(G) = \Theta(r^{-2})$ for $G \in \mathcal{G}(n, r, 1)$, the usual random geometric graph, provided that $\pi nr^2 - \ln n \rightarrow \infty$ (that is, above the connectivity threshold). For the percolated random geometric graph $G \in \mathcal{G}(n, r, p)$, we show that with high probability $\mathcal{AC}(G) = \Theta(r^{-2}p^{-1} \ln n)$, provided that $pnr^2 \geq n^{1/2+\varepsilon}$ and $p < 1 - \varepsilon$ for some $\varepsilon > 0$.

© 2015 Elsevier Ltd. All rights reserved.

1. Introduction and statement of results

In this paper, we study the following graph process, which was recently introduced by Benjamini, Shinkar, and Tsur [4]. Let $G = (V, E)$ be a finite connected graph. We start the process by placing exactly one *agent* on each vertex of G . Every pair of agents on adjacent vertices is declared to be *acquainted*, and remains so throughout the process. In each round of the process, we choose some matching M in G . (M need not be maximal; perhaps it is a single edge.) For each edge of M , we swap the agents occupying its endpoints, which may cause more agents to become acquainted. The

E-mail addresses: t.muller@uu.nl (T. Müller), pralat@ryerson.ca (P. Prałat).

acquaintance time of G , denoted by $\mathcal{AC}(G)$, is the minimum number of rounds required for all agents to become acquainted with one another.

It is clear that

$$\mathcal{AC}(G) \geq \frac{\binom{|V|}{2}}{|E|} - 1, \quad (1)$$

since $|E|$ pairs are acquainted initially, and at most $|E|$ new pairs become acquainted in each round. In [4], it was shown that always $\mathcal{AC}(G) = O(\frac{n^2}{\ln n / \ln \ln n})$, where $n = |V|$, which was slightly sharpened in [17] to $\mathcal{AC}(G) = O(\frac{n^2}{\ln n})$. This general upper bound was recently improved and now we know that $\mathcal{AC}(G) = O(n^{3/2})$ for every graph G , which was conjectured in [4] and is tight up to a multiplicative constant [2]. In [17], another conjecture from [4] on the acquaintance time of the random graph $\mathcal{G}(n, p)$ was investigated. In [4], it was conjectured that asymptotically almost surely $\mathcal{AC}(G(n, p)) \leq \ln^{O(1)} n/p$, provided that $G(n, p)$ is connected. In [17], it was shown that asymptotically almost surely $\mathcal{AC}(G(n, p)) = O(\ln n/p)$, provided that $pn - \ln n - \ln \ln n \rightarrow \infty$ as $n \rightarrow \infty$ (that is, above the threshold for Hamiltonicity). Moreover, a matching lower bound for dense random graphs was provided, which also implies that asymptotically almost surely K_n cannot be covered with $o(\ln n/p)$ copies of a random graph $\mathcal{G}(n, p)$, provided that $pn > n^{1/2+\varepsilon}$ and $p < 1 - \varepsilon$ for some $\varepsilon > 0$. In [10], the task was finished by showing that the conjecture holds in the strongest possible sense, that is, it holds right at the time the random graph process creates a connected graph.

The problem is similar in flavour to Gossiping and Broadcasting [13], Collision-Free Network Exploration [9] and Target Set Selection [15,7,26]. One problem of particular relevance is Routing Permutations on Graphs via Matchings [1].

In the present paper, we consider the acquaintance time of (percolated) random geometric graphs. If $V \subseteq \mathbb{R}^2$ is a set of points and $r > 0$ then the *geometric graph* $\mathcal{G}(V, r)$ is the graph with vertex set V and an edge between two points if and only if their distance is at most r . Such a graph is also called a *unit disk graph* since it is the intersection graph of disks of the same radius (namely disks of radius $r/2$ centred on the points of V). Throughout this paper, we let $X_1, X_2, \dots \in \mathbb{R}^2$ be an infinite supply of random points, i.i.d. on the unit square. For notational convenience (and following Penrose [22]) we set:

$$\mathcal{X}_n := \{X_1, X_2, \dots, X_n\}. \quad (2)$$

The *random geometric graph* $\mathcal{G}(n, r)$ is the random graph obtained by taking \mathcal{X}_n as the vertex set, i.e. $\mathcal{G}(n, r) := \mathcal{G}(\mathcal{X}_n, r)$. To prevent dealing with annoying trivial cases we shall always assume that $r < \sqrt{2}$ throughout this paper (otherwise $G \in \mathcal{G}(n, r)$ is a clique and $\mathcal{AC}(G) = 0$).

The study of random geometric graphs essentially goes back to Gilbert [11] who defined a very similar model in 1961. For this reason it is often also called the *Gilbert model*. Random geometric graphs have been the subject of a considerable research effort in the last two decades. As a result, detailed information is now known on various aspects such as (k -)connectivity [23,24], the largest component [22], the chromatic number and clique number [21,19], the (non-)existence of Hamilton cycles [3,20] and the simple random walk on the graph [8]. A good overview of the results prior to 2003 can be found in the monograph [22].

The *percolated random geometric graph* $\mathcal{G}(n, r, p)$ is obtained by retaining each edge of $\mathcal{G}(n, r)$ with probability p (and discarding it with probability $1 - p$). To be more precise, for each edge of $\mathcal{G}(n, r)$ we flip a biased coin which is independent of \mathcal{X}_n and the other coin tosses for the other edges, and keep the edge if the coin comes up heads. In particular, $\mathcal{G}(n, r) = \mathcal{G}(n, r, 1)$. This model has not received the same amount of attention from mathematicians as the usual random geometric graph, but it is the obvious next special case to consider of the *soft random geometric graph*. In a soft random geometric graph a pair of vertices is connected with a probability that is a function of their distance (both the usual and the percolated random geometric graph obviously fit this framework). Soft random geometric graphs can count on interest from the engineering community (see for instance [12,18,27]), as they are considered more realistic models of wireless networks than usual random geometric graphs. Recently, Penrose [25] gave a precise result on the threshold for connectivity in soft random geometric graphs.

As typical in random graph theory, we shall consider only asymptotic properties of $\mathcal{G}(n, r)$ and $\mathcal{G}(n, r, p)$ as $n \rightarrow \infty$, where both r and p may and usually do depend on n . Throughout this paper, we will say that a sequence of events E_1, E_2, \dots holds *with high probability* (abbreviated w.h.p.) if $\mathbb{P}(E_n) \rightarrow 1$ as $n \rightarrow \infty$.

It follows from a very precise result of Penrose [23] that the (classical) random geometric graph $\mathcal{G}(n, r_n)$ is w.h.p. connected if and only if the sequence $(r_n)_n$ is such that $\pi n r_n^2 - \ln n \rightarrow \infty$ as $n \rightarrow \infty$. We are able to obtain the following result, which tells us the likely value of the acquaintance time up to a constant factor, whenever the acquaintance time is (w.h.p.) well defined.

Theorem 1.1. *If $(r_n)_n$ is such that $\pi n r_n^2 - \ln n \rightarrow \infty$, then $\mathcal{AC}(G(n, r_n)) = \Theta(r_n^{-2})$ w.h.p.*

For the percolated random geometric graph $\mathcal{G}(n, r_n, p_n)$ we are slightly less successful. For dense graphs we determine the likely value of the acquaintance time up to a multiplicative constant, but the behaviour for sparser graphs remains undetermined.

Theorem 1.2. *Let $\varepsilon > 0$ be arbitrary. If $(r_n)_n$ and $(p_n)_n$ are such that $p_n < 1 - \varepsilon$ and $p_n n r_n^2 \geq n^{1/2+\varepsilon}$, then $\mathcal{AC}(\mathcal{G}(n, r_n, p_n)) = \Theta(r_n^{-2} p_n^{-1} \ln n)$ w.h.p.*

In the course of the proof we will in fact prove slightly more. Namely, we will derive an upper bound of $\mathcal{AC}(G) = O(r_n^{-2} p_n^{-1} \ln n)$ that works whenever $p_n n r_n^2 \geq K \ln n$ for some large constant K .

2. Preliminaries

Throughout this paper $B(x, r) \subseteq \mathbb{R}^2$ will denote the points at distance $< r$ to the point x . We will denote by $\text{Po}(\lambda)$ the Poisson distribution with parameter λ , and $\text{Bi}(n, p)$ will denote the binomial distribution with parameters n and p . We will make use of the following incarnation of the Chernoff bounds. A proof can, for instance, be found in Chapter 1 of [22].

Lemma 2.1. *Let Z be either Poisson or Binomially distributed, and write $\mu := \mathbb{E}Z$.*

(i) *For all $k \geq \mu$ we have*

$$\mathbb{P}(Z \geq k) \leq e^{-\mu H(k/\mu)},$$

(ii) *For all $k \leq \mu$ we have*

$$\mathbb{P}(Z \leq k) \leq e^{-\mu H(k/\mu)},$$

where $H(x) := x \ln x - x + 1$.

We will need the following standard and elementary result. A proof can be found in [5].

Lemma 2.2. *Let G be a connected geometric graph. Then G has a spanning tree of maximum degree at most five.*

We remark that the number five in this last lemma is only true in dimension two (but in every dimension a similar lemma can be proved with another dimension-dependent constant).

We will make use of some known results on the acquaintance time of graphs. The first one is a very recent result of Angel and Shinkar [2].

Theorem 2.3 ([2]). *For every connected graph G we have $\mathcal{AC}(G) \leq 20 \cdot \Delta(G) \cdot |V(G)|$, where $\Delta(G)$ is the maximum degree of G .*

Recall that $G[H]$ is the graph with vertex set $V(G) \times V(H)$ with an edge between (u, v) and (u', v') if either (1) $u = u'$ and $vv' \in E(H)$, or (2) $uu' \in E(G)$. So, in particular, $G[K_s]$ is the graph we get by replacing each vertex of G by an s -clique and adding all edges between the cliques corresponding to adjacent vertices of G . We will use the following two straightforward observations.

Lemma 2.4. We have $\mathcal{AC}(G[K_s]) \leq \mathcal{AC}(G)$ for all connected G and all $s \in \mathbb{N}$.

Sketch of the Proof. We partition the agents into groups of size s , corresponding to the s -cliques that have replaced the vertices of G . Each group is treated as a single agent, and the strategy that yields $\mathcal{AC}(G)$ is used. \square

Lemma 2.5. Let G be an arbitrary connected graph. There is a strategy such that all agents get acquainted, and return to their initial vertices in $2 \cdot \mathcal{AC}(G)$ rounds.

Sketch of the Proof. Follow the strategy that yields $\mathcal{AC}(G)$. Then, simply repeat the sequence of moves in reversed order. \square

We will also use the fact, observed in [4], that for any graph G on n vertices with a Hamiltonian path, we have $\mathcal{AC}(G) = O(n)$. In fact, we need a slightly stronger statement that was proved in [17].

Lemma 2.6 ([17]). Let G be a graph on n vertices. If G has a Hamiltonian path, then there exists a strategy ensuring that within $2n$ rounds every pair of agents gets acquainted and, moreover, that every agent visits every vertex.

3. Proof of the lower bound in Theorem 1.1

The following Lemma is a simplification of Lemma A.1 in [21], where slightly more is proved.

Lemma 3.1 ([21]). If the sequence $(r_n)_n$ is such that $n^2 r_n^2 \rightarrow \infty$ then the number of edges of $\mathcal{G}(n, r_n)$ is $\Theta(n^2 r_n^2)$ w.h.p.

The lower bound on \mathcal{AC} now follows immediately from the trivial lower bound (1). We see that when $\pi n r_n^2 = \ln n + \omega(1)$ then we have

$$\mathcal{AC}(G) \geq \frac{\binom{n}{2}}{|E(G)|} - 1 = \Omega(r_n^{-2}) \quad \text{w.h.p.}, \quad (3)$$

and the proof of the lower bound is finished.

4. The proof of the upper bound in Theorem 1.1

It is convenient to split the proof into two cases. The first, and easier, case is when the sequence $(r_n)_n$ is such that $\pi n r_n^2$ is at least $K \ln n$ for some large constant $K > 0$. We can then make use of the concentration phenomenon to give a relatively easy proof of an upper bound of the right order of magnitude. The second, and more involved, case is when $(r_n)_n$ is such that $\pi n r_n^2$ is somewhere between $\ln n + \omega(1)$ and $K \ln n$. Here, we need to make use the detailed information on the structure of $G(n, r_n)$ close to the connectivity threshold. Luckily, a lot of this structural information has previously been obtained in, for instance, [5], and we can obtain the statement we need for our proofs by adapting some previous results to suit our needs. We shall refer to the first case as the dense case, and to the second as the sparse case.

4.1. The upper bound for $G(n, r_n)$ in the dense case

Throughout the rest of the paper, we let $\mathcal{D}(m)$ denote the dissection of $[0, 1]^2$ into squares of side length $s(m) := 1/m$, with $m \in \mathbb{N}$. We will call the squares of this dissection *cells*, and given cell $c \in \mathcal{D}(m)$, we will denote by $V(c)$ the set of points of \mathcal{X}_n that fall in c .

For $(r_n)_n$, an arbitrary sequence of numbers with $0 < r_n < \sqrt{2}$, let us define $m_n := \lceil 1000/r_n \rceil$. Then we have $1/m_n \leq r_n/1000$ and $1/m_n = \Omega(r_n)$. (Moreover $1/m_n \sim r_n/1000$ if $r_n \rightarrow 0$.) Let us write $\mathcal{D}_n := \mathcal{D}(m_n)$ and let $\mu_n := n/m_n^2$ denote the expectation $\mathbb{E}|V(c)|$.

Lemma 4.1. *There is a constant $K > 0$ such that if $(r_n)_n$ is such that $\pi n r_n^2 \geq K \ln n$ then $0.9 \cdot \mu_n \leq |V(c)| \leq 1.1 \cdot \mu_n$ for all $c \in \mathcal{D}_n$, w.h.p.*

Proof. Fix a cell $c \in \mathcal{D}_n$. By Lemma 2.1 the probability that $|V(c)| > 1.1 \cdot \mu_n$ satisfies

$$\mathbb{P}(|V(c)| > 1.1 \cdot \mu_n) \leq e^{-\mu_n H(1.1)},$$

where $H(x) = x \ln x - x + 1$ is as in Lemma 2.1. Now notice that m_n is non-decreasing in r_n , so that $\mu_n = n/m_n^2$ is non-increasing in r_n . It follows that whenever $\pi n r_n^2 \geq K \ln n$ for some constant $K > 0$ then

$$\mathbb{P}(|V(c)| > 1.1 \cdot \mu_n) \leq \exp\left[-(1 + o(1)) \cdot \frac{H(1.1)K}{\pi 10^6} \cdot \ln n\right].$$

(Here we have used that if $\pi n r_n^2 = K \ln n$ then $1/m_n \sim r_n/1000$.) Hence, by the union bound,

$$\begin{aligned} \mathbb{P}(\text{There exists a } c \in \Gamma_n \text{ with } |V(c)| > 1.1 \cdot \mu_n) &\leq m_n^2 \cdot \exp\left[-(1 + o(1)) \cdot \frac{H(1.1)K}{\pi 10^6} \cdot \ln n\right] \\ &\leq n^2 \cdot n^{-(1+o(1))H(1.1)K\pi^{-1}10^{-6}} \\ &= o(1), \end{aligned}$$

provided K is chosen sufficiently large. Completely analogously, we can show that, w.h.p., no cell will have less than $0.9 \cdot \mu_n$ points, provided we chose K sufficiently large. \square

For the remainder of the proof, let $V \subseteq [0, 1]^2$, $0 < r < \sqrt{2}$ be such that the conclusion of this last lemma holds, but otherwise arbitrary. It suffices to show that $G := G(V, r)$ satisfies $\mathcal{AC}(G) = O(r^{-2})$.

For each $c \in \mathcal{D}_n$ we partition $V(c)$ into three parts $V_1(c), V_2(c), V_3(c)$, each of cardinality at most $0.4 \cdot \mu_n$. For each pair $1 \leq i < j \leq 3$ and each cell $c \in \Gamma_n$, let $W_{ij}(c) \subseteq V(c)$ be a set of cardinality exactly $t := \lfloor 0.9\mu_n \rfloor$ such that $V_i(c) \cup V_j(c) \subseteq W_{ij}(c)$; and set $W_{ij} := \bigcup_{c \in \Gamma_n} W_{ij}(c)$. Let $G_{ij} = G[W_{ij}]$ denote the subgraph induced by W_{ij} . We now observe that, since points in touching cells of the dissection \mathcal{D}_n have distance at most r , the graph G_{ij} has a spanning subgraph that is isomorphic to $H[K_t]$ where H denotes the $m_n \times m_n$ -grid. It follows from Theorem 2.3 and Lemmas 2.4 and 2.5 that we can acquaint all agents on vertices of W_{ij} with each other, and return them to their starting positions in $O(m_n^2) = O(r_n^{-2})$ rounds.

By repeating this procedure for each of W_{12}, W_{13}, W_{23} , we acquaint all agents with each other in $O(r_n^{-2})$ rounds, as required.

4.2. Structural definitions and lemmas needed for the sparse case

Before we can start the last part of the proof of Theorem 1.1, we need to recall some definitions and (slightly adapted versions of) results from [5]. Here we simply list the definitions we need, without further discussion or motivation. For a more gentle introduction to the definitions, including a diagram, we refer the reader to [5] (Section 1.1 and the beginning of Section 3). Let us consider any geometric graph $G = (V, r)$, where $V = \{x_1, x_2, \dots, x_n\} \subset [0, 1]^2$.

Let $m \in \mathbb{N}$ be such that $s(m) := 1/m \leq r/1000$ and consider $\mathcal{D} = \mathcal{D}(m)$. Given $T > 0$ and $V \subseteq [0, 1]^2$, we call a cell $c \in \mathcal{D}$ *good* with respect to T, V if $|c \cap V| \geq T$ and *bad* otherwise. When the choice of T and V is clear from the context we will just speak of good and bad. Let $\Gamma = \Gamma(V, m, T, r)$ denote the graph whose vertices are the good cells of $\mathcal{D}(m)$, with an edge $cc' \in E(\Gamma)$ if and only if the lower left corners of c, c' have distance at most $r - s\sqrt{2}$. (Note that this way, any $x \in c$ and $y \in c'$ have distance $\|x - y\| \leq r$.) We will usually just write Γ when the choice of V, m, T, r is clear from the context. Let us denote the components of Γ by $\Gamma_1, \Gamma_2, \dots$ where Γ_i has at least as many cells as Γ_{i+1} (ties are broken arbitrarily). For convenience we will also write $\Gamma_{\max} = \Gamma_1$. We will often be a bit sloppy and identify Γ_i with the union of its cells, and speak of $\text{diam}(\Gamma_i)$ and the distance between Γ_i and Γ_j and so forth.

Let us call a point $v \in V$ *safe* if there is a good cell $c \in \Gamma_{\max}$ such that $|B(v; r) \cap V \cap c| \geq T$. (That is, in the geometric graph $G(V; r)$, the point v has at least T neighbours inside c .) Otherwise, if there

is a good cell $c \in \Gamma_i$, $i \geq 2$, such that $|B(v; r) \cap V \cap c| \geq T$, we say that v is *risky*. Otherwise we call v *dangerous*.

For $i \geq 2$ we let Γ_i^+ denote the set of all points of V in cells of Γ_i , together with all risky points v that satisfy $|B(v; r) \cap V \cap c| \geq T$ for at least one $c \in \Gamma_i$. The following is a list of desirable properties that we would like V and $\Gamma(V, m, T, r)$ to have:

- (str-1) Γ_{\max} contains more than $0.99 \cdot |\mathcal{D}|$ cells;
- (str-2) $\text{diam}(\Gamma_i^+) < r/100$ for all $i \geq 2$;
- (str-3) If $u, v \in V$ are dangerous then either $\|u - v\| < r/100$ or $\|u - v\| > r \cdot 10^{10}$;
- (str-4) For all $i > j \geq 2$ the distance between Γ_i^+ and Γ_j^+ is at least $r \cdot 10^{10}$;
- (str-5) If $v \in V$ is dangerous and $i \geq 2$ then the distance between v and Γ_i^+ is at least $r \cdot 10^{10}$.

Finally, we introduce some terminology for sets of dangerous and risky points. Suppose that $V \subseteq [0, 1]^2$ and m, T, r are such that (str-1)–(str-5) above hold. Dangerous points come in groups of points of diameter $< r/100$ that are far apart. We formally define a *dangerous cluster* (with respect to V, m, T, r) to be an inclusion-wise maximal subset of V with the property that $\text{diam}(A) < r \cdot 10^{10}$ and all elements of A are dangerous.

A set $A \subseteq V$ is an *obstruction* (with respect to V, m, T, r) if it is either a dangerous cluster or Γ_i^+ for some $i \geq 2$. We call A an *s-obstruction* if $|A| = s$. By (str-3)–(str-5), obstructions are pairwise separated by distance $r \cdot 10^{10}$. (One consequence: a vertex in a good cell is adjacent in G to at most one obstruction.) A point $v \in V$ is *crucial* for A if

- (cruc-1) $A \subseteq N(v)$, and;
- (cruc-2) v is safe.

We are interested in the following choice of m for our dissection. For $n \in \mathbb{N}$ and $\eta > 0$ a constant, let us define

$$m_n := \left\lceil \sqrt{\frac{n}{\eta^2 \ln n}} \right\rceil. \quad (4)$$

The following lemma is almost identical to a lemma in [5]. For completeness we spell out the adaptations that need to be made to its proof in [Appendix A](#).

Lemma 4.2. *For every sufficiently small $\eta > 0$, there exists a $\delta = \delta(\eta) > 0$ such that the following holds. Let m_n be given by (4), let \mathcal{X}_n be as in (2), let $T_n \leq \delta \ln n$ and let r_n be such that $\pi n r_n^2 = \ln n + o(\ln n)$. Then (str-1)–(str-5) hold for $\Gamma(\mathcal{X}_n, m_n, T_n, r_n)$ w.h.p.*

The following lemma is also a slightly adapted version of a result in [5]. Its proof can be found in [Appendix B](#).

Lemma 4.3. *For every sufficiently small $\eta > 0$, there exists a $\delta = \delta(\eta) > 0$ such that the following holds. Let $(m_n)_n$ be given by (4), let $T_n \leq \delta \ln n$ and let $V_n := \mathcal{X}_n$ with \mathcal{X}_n as in (2), let $(r_n)_n$ be a sequence of positive numbers such that $\pi r_n^2 - \ln n \rightarrow \infty$.*

Then, w.h.p., it holds that for every $s \geq 2$, every s -obstruction has at least $s - 100$ crucial vertices.

The proof of the following lemma is analogous to that of [Lemma 4.1](#) and is left to the reader.

Lemma 4.4. *If $\eta > 0$ is fixed, $(m_n)_n$ is as given by (4) and $V_n := \mathcal{X}_n$ then there is a constant C such that, w.h.p., every cell contains at most $C \ln n$ points.*

4.3. The proof of the upper bound in Theorem 1.1 in the sparse case

It remains to prove the upper bound of Theorem 1.1 for a sequence r_n such that $\pi n r_n^2 = \ln n + \omega(1)$ and $\pi n r_n^2 \leq K \ln n$ for a large constant K . For this range it suffices to consider the case when $1 \ll \pi n r_n^2 - \ln n \ll \ln n$ and prove that in that case w.h.p. $\mathcal{AC}(\mathcal{G}(n, r_n)) = O(n/\ln n)$. (By (3) we already have the asymptotically almost sure lower bound $\mathcal{AC}(\mathcal{G}(n, r_n)) = \Omega(r_n^{-2}) = \Omega(n/\ln n)$ for all sequences $(r_n)_n$ satisfying $\pi n r_n^2 \leq K \ln n$.) Let us thus pick such a sequence r_n and assume η, δ etc. have been chosen in such a way that the conclusions of Lemmas 4.2 and 4.3 hold w.h.p. We also know from [23] that in this range $\mathcal{G}(n, r_n)$ is connected w.h.p.

In the sequel of the proof we let $V \subseteq [0, 1]^2$ be an arbitrary set of points, and $r, m, \eta, \delta > 0$ be arbitrary numbers such that $G(V, r)$ is connected and the conclusions of Lemmas 4.2–4.4 are satisfied. It suffices to show that every such graph $G = G(V, r)$ has acquaintance time $O(n/\ln n)$, as we will now show.

Claim 4.5. *There is a constant c_1 such that every obstruction consists of at most $c_1 \ln n$ points.*

Proof. Every obstruction O has diameter $r/100$ and hence there are only $O(1)$ cells that contain points of O . Since each cell contains $O(\ln n)$ points, we are done. \square

Each point v that is safe, but not in a cell of Γ_{\max} has at least T neighbours in some cell $c \in \Gamma_{\max}$. We arbitrarily “assign” v to such a cell.

Claim 4.6. *There is a constant $c_2 > 1$ such that the following holds. For every obstruction O there is a cell $c \in \Gamma_{\max}$ such that at least $|O|/c_2$ vertices that are crucial for O have been assigned to c .*

Proof. Since G is connected, every obstruction has at least one crucial vertex. (It is adjacent to at least one vertex $v \in V \setminus O$ and this v cannot be dangerous or risky.) This shows that, by choosing c_2 sufficiently large, the claim holds whenever $|O| < 1000$. Let us thus assume $|O| \geq 1000$. Since O has diameter $< r/100$, there is a constant D such that the crucial vertices are all assigned to all one of the D cells within range $2r$ of O . Since $|O| > 1000$ there are at least $|O| - 100 > |O|/2$ crucial vertices for O , and hence at least $|O|/2D$ of these crucial vertices are assigned to the same cell. \square

For each obstruction O , we now assign all its vertices to a cell $c \in \Gamma_{\max}$ such that at least $|O|/c_2$ vertices that are crucial for O have been assigned to c . For a cell $c \in \Gamma$ let $V(c)$ denote the set of points that fell in c . For each $c \in \Gamma_{\max}$, let $A(c)$ denote the union of $V(c)$ with all vertices that have been assigned to c . Let us remark that:

Claim 4.7. *There exists a constant c_3 such that $|A(c)| \leq c_3 \ln n$ for all $c \in \Gamma_{\max}$.*

Proof. Note that if a vertex v has been assigned to c , it must lie within distance $2r$ of c . Hence there are only $O(1)$ cells in which $A(c)$ is contained, and since every cell contains $O(\ln n)$ points, we are done. \square

Let us now partition $A(c)$ into sets $A_1(c), A_2(c), \dots, A_L(c)$ each of size at most $T/100$, with $L = \lceil c_3/(100\delta) \rceil$. The following observation will be key to our strategy.

Claim 4.8. *There is a constant c_4 such that the following holds. For every $c \in \Gamma_{\max}$ and every $A' \subseteq A(c)$ with $|A'| \leq T/50$ there is a sequence of at most c_4 moves that results in the agents on vertices on A' being placed on vertices of $V(c)$ (and uses only edges of $G[A(c)]$).*

Proof. We first move all agents on vertices of $A' \setminus V(c)$ on safe vertices not in $V(c)$ onto vertices of $V(c) \setminus A'$ in one round. (To see that this can be done, note that $|A'| < T/50$ and each safe vertex of A' has at least T neighbours in $V(c)$.) Let $W \subseteq V(c)$ be the set of vertices now occupied by agents that were originally on A' .

If A' also contains (part of) some obstruction O , then we partition $O \cap A'$ into $O(1)$ sets O_1, O_2, \dots, O_K of cardinality at most $|O|/c_2$ where $K \leq \lceil 1/c_2 \rceil$ (and hence is a constant). We first move the agents on vertices O_1 onto crucial vertices assigned to c , and then on vertices of $V(c) \setminus W$, in two rounds. (Note

this is possible since $|A'| \leq T/50$ and each crucial vertex is adjacent to at least T vertices of $V(c)$.) Similarly, supposing that the agents on O_1, O_2, \dots, O_{i-1} have already been moved onto vertices of $V(c)$, we can move the agents on vertices of O_i onto vertices of $V(c) \setminus W$ not occupied by agents from O_1, O_2, \dots, O_{i-1} in two rounds.

We thus have moved all agents on vertices of A' onto vertices of $V(c)$ in a constant number of rounds, as required. \square

We are now ready to describe the overall strategy. Let us write $A_i := \bigcup_{c \in \Gamma_{\max}} A_i(c)$. For each pair of indices $1 \leq i < j \leq L$ we do the following.

First we move all agents of $A_i(c) \cup A_j(c)$ onto vertices of $V(c)$ (in constantly many moves, simultaneously for all cells $c \in \Gamma_{\max}$). Next, we select a set $B(c) \subseteq V(c)$ for each $c \in \Gamma_{\max}$ with $|B(c)| = T$ and all agents that were on $A_i(c) \cup A_j(c)$ originally are now on vertices of $B(c)$. By Lemma 2.2, the largest component Γ_{\max} of the cells-graph has a spanning tree H of maximum degree at most five. Thus, by Theorem 2.3, we have $\mathcal{AC}(\Gamma_{\max}) \leq O(\Delta(H) \cdot |V(H)|) = O(|\Gamma|) = O(n/\ln n)$. Now note that the graph spanned by $\bigcup_{c \in \Gamma_{\max}} B(c)$ contains a spanning subgraph isomorphic to $\Gamma_{\max}[K_T]$. Using Lemmas 2.4 and 2.5, we can thus acquaint all vertices of $\bigcup_{c \in \Gamma_{\max}} B(c)$ with each other and return them to their starting vertices in $O(n/\ln n)$ rounds. So, in particular, we have acquainted the agents of $A_i \cup A_j$ and returned them to their starting positions, in $O(n/\ln n)$ moves.

Once we have repeated this procedure for each of the $\binom{L}{2} = O(1)$ pairs of indices all agents will be acquainted, still in $O(n/\ln n)$ rounds. This concludes the proof of Theorem 1.1.

5. The proof of the lower bound in Theorem 1.2

In hopes of doing better than the trivial lower bound (1) on the acquaintance time of $\mathcal{G}(n, r_n, p)$, we consider a variant of the original process. This approach was used for binomial random graphs [17] and, after some adjustments combined with an additional averaging argument, can be used here as well. Suppose that each agent has a helicopter and can, on each round, move to any vertex she wants. (We retain the requirement that no two agents can occupy a single vertex simultaneously.) In other words, in every step of the process, the agents choose some permutation π of the vertices, and the agent occupying vertex v flies directly to vertex $\pi(v)$, regardless of whether there is an edge or even a path between v and $\pi(v)$. (In fact, it is no longer necessary that the graph be connected.) Let the *helicopter acquaintance time* $\overline{\mathcal{AC}}(G)$ be the counterpart of $\mathcal{AC}(G)$ under this new model, that is, the minimum number of rounds required for all agents to become acquainted with one another. Since helicopters make it easier for agents to get acquainted, we immediately get that for every graph G ,

$$\overline{\mathcal{AC}}(G) \leq \mathcal{AC}(G). \quad (5)$$

On the other hand, $\overline{\mathcal{AC}}(G)$ also represents the minimum number of copies of a graph G needed to cover all edges of a complete graph of the same order. Thus inequality (1) can be strengthened to $\overline{\mathcal{AC}}(G) \geq \binom{|V|}{2} / |E| - 1$.

In order to prove the lower bound in part (ii) of Theorem 1.2, we prove the following general result. If G is a graph then we denote by G^p the random subgraph of G in which every edge is kept with probability p and discarded with probability $1 - p$ (independently of all other edges). So, in particular, K_n^p is the familiar binomial random graph $G(n, p)$ and $\mathcal{G}(n, r_n, p)$ is the same as $\mathcal{G}^p(n, r_n)$.

Theorem 5.1. *Let $\varepsilon > 0$ be arbitrary. If $(G_n)_n$ is a sequence of graphs with $v(G_n) = n$, and $(p_n)_n$ is a sequence of edge-probabilities satisfying $p_n \leq 1 - \varepsilon$, and $p_n e(G_n) \geq n^{3/2+\varepsilon}$ for all n , then*

$$\overline{\mathcal{AC}}(G_n^{p_n}) = \Omega\left(\frac{n^2 \ln n}{p_n \cdot e(G_n)}\right) \quad \text{w.h.p.}$$

Proof. Let a_1, a_2, \dots, a_n denote the n agents, and let $A = \{a_1, a_2, \dots, a_n\}$. Take

$$k = \frac{\varepsilon}{20} (n^2 / e(G_n)) \cdot \log_{1/(1-p_n)} n = \Theta\left(\frac{n^2 \ln n}{p_n \cdot e(G_n)}\right)$$

and fix k bijections $\pi_i : A \rightarrow V(G_n)$, for $i \in \{0, 1, \dots, k-1\}$. This corresponds to fixing a $(k-1)$ -round strategy for the agents; in particular, agent a_j occupies vertex $\pi_i(a_j)$ in round i . We aim to show that at the end of the process (that is, after $k-1$ rounds) the probability that all agents are acquainted is only $o((1/n!)^k)$. This will complete the proof. Indeed, the number of choices for $\pi_0, \pi_1, \dots, \pi_{k-1}$ is $(n!)^k$, so by the union bound, w.h.p. no strategy makes all pairs of agents acquainted.

We say that a pair of vertices (or agents) is *reachable* in a particular round if they are on vertices that are adjacent in the non-percolated graph G_n . (So a pair of agents can be reachable during one round but not reachable during another one.) Let b be the number of pairs of agents that are reachable during more than $10kn^{-2}e(G_n)$ rounds (with respect to the given agents' strategy, that is, the k bijections that are fixed). In each round, at most $e(G_n)$ pairs are reachable. It follows that $b \cdot 10kn^{-2}e(G_n) \leq e(G_n)k$, so that

$$b \leq n^2/10 \leq \binom{n}{2}/2.$$

Hence, at least half of all pairs of agents are reachable during at most $10ke(G_n)/n^2$ rounds. We call these pairs of agents *important*.

To estimate the probability that a given agents' strategy makes all pairs of important agents acquainted, we consider the following analysis, which iteratively exposes edges of a percolated random graph $G_n^{p_n}$. For any pair $q = \{a_x, a_y\}$ of important agents, we consider all reachable pairs of vertices visited by this pair of agents throughout the process:

$$S(q) = \{e \in E(G_n) : e = \pi_i(a_x)\pi_i(a_y) \text{ for some } i \in \{0, 1, \dots, k-1\}\}.$$

Since q is important, $1 \leq |S(q)| \leq 10ke(G_n)/n^2$. Let us now fix an arbitrary ordering q_1, q_2, \dots, q_m of our important pairs and consider the following process. We take the first pair q_1 of important agents and expose the edges of $G_n^{p_n}$ in $S(q_1)$, one by one until we either find an edge that is present in $G_n^{p_n}$ or we have exposed all of $S(q_1)$. If we expose all of $S(q_1)$ without discovering an edge, then the pair q_1 never gets acquainted and we halt our procedure. If instead we do discover some edge e of $G_n^{p_n}$, then we discard all pairs of important agents that ever occupy this edge (that is, we discard all pairs q such that $e \in S(q)$). We now shift our attention to the next pair q_i of important agents that we did not yet discard and repeat the procedure of exposing edges $S(q_i)$ until we find one that acquaints q_i or we run out of edges. It may happen that some of the pairs of vertices in $S(q_i)$ have already been exposed, but the analysis guarantees that no edge has yet been discovered.

We continue this process until either we have found an important pair that never gets acquainted or all available pairs of important agents have been investigated. Considering one pair of important agents can force us to discard at most k important pairs (including the original pair) since in each round the edge acquaints at most one pair. Hence, the process investigates at least $\frac{1}{2} \binom{n}{2}/k$ pairs of important agents. Moreover, writing $E_t := \{q_1, q_2, \dots, q_{t-1} \text{ get acquainted and } q_t \text{ did not get discarded}\}$, we have

$$\begin{aligned} \mathbb{P}(q_t \text{ gets acquainted} | E_t) &\leq 1 - (1-p)^{|S(q_t)|} \\ &\leq 1 - (1-p)^{10kn^{-2}e(G_n)}. \end{aligned}$$

Hence, we find

$$\begin{aligned} \mathbb{P}(\text{all pairs acquainted}) &\leq \mathbb{P}(\text{all important pairs acquainted}) \\ &\leq \left(1 - (1-p)^{10kn^{-2}e(G_n)}\right)^{\frac{1}{2} \binom{n}{2}/k} \\ &\leq \exp \left[-(1-p)^{10kn^{-2}e(G_n)} \cdot \frac{1}{2} \binom{n}{2} / k \right] \\ &\leq \exp \left[-n^{-\varepsilon/2} \cdot \frac{1}{2} \binom{n}{2} / k \right] \\ &= \exp \left[-\Omega(n^{-\varepsilon/2} p_n e(G_n) / \ln n) \right], \end{aligned}$$

using that $k = \frac{\varepsilon}{20} (n^2/e(G_n)) \cdot \log_{1/(1-p_n)} n$ for the third line and that $k = \Theta(n^2 \ln n / (p_n e(G_n)))$ for the last line. Now note that

$$\begin{aligned} (n!)^k &\leq n^{k \cdot n} \\ &= \exp[k \cdot n \ln n] \\ &= \exp\left[O\left(n^3 \ln^2 n / (p_n e(G_n))\right)\right]. \end{aligned}$$

Since $p_n e(G_n) \geq n^{1/2+\varepsilon}$ we also have that $n^{-\varepsilon/2} p_n e(G_n) / \ln n \gg n^3 \ln^2 n / (p_n e(G_n))$, and hence

$$\mathbb{P}(\text{all pairs acquainted}) = o((1/n!)^k),$$

which concludes the proof by a previous remark. \square

Combining the last theorem with [Lemma 3.1](#), we immediately get:

Corollary 5.2. *Let $\varepsilon > 0$ be arbitrary. If the sequences $(r_n)_n$ and $(p_n)_n$ are such that $p_n < 1 - \varepsilon$ for all n and $p_n^2 n r_n^2 \geq n^{1/2+\varepsilon}$, then*

$$\mathcal{AC}(\mathcal{G}(n, r_n, p_n)) \geq \overline{\mathcal{AC}}(\mathcal{G}(n, r_n, p_n)) = \Omega(r_n^{-2} p_n^{-1} \ln n) \quad \text{w.h.p.}$$

6. The proof of the upper bound in [Theorem 1.2](#)

Let us start with the following useful observation. Let $p = p(n)$ and $t = t(n)$, and let $\mathcal{B}(t, p)$ be the “standard” random bipartite graph with bipartite sets X and Y such that $|X| = |Y| = t$. For each pair of vertices $x \in X$ and $y \in Y$, we introduce an edge xy with probability p , independently of all other edges. We will consider the probability that $\mathcal{B}(t, p)$ has a perfect matching. Very precise information is already known about perfect matchings in this random graph model (see for instance [\[6\]](#), Section 7.3). We however need precise quantitative bounds on the probability of existence of a perfect matching, which do not appear in the literature as far as we are aware of.

Lemma 6.1. *With $\mathcal{B}(t, p)$ the random bipartite graph as above, we have*

$$\mathbb{P}(\mathcal{B}(t, p) \text{ has a perfect matching}) = 1 - O(t \cdot e^{-\gamma t p}),$$

for some universal constant $\gamma > 0$.

Proof. Let us first observe that if $tp \leq K \ln t$ for some constant K then there is nothing to prove since we may assume, without loss on generality, that $\gamma > 0$ is sufficiently small for $t \cdot e^{-\gamma t p} \rightarrow \infty$ to hold in this case. Let us thus assume that $tp > K \ln t$ in the sequel, where $K > 0$ is a constant to be chosen more precisely later on in the proof.

Set $s_0 = \max\{s \in \mathbb{N} : ps \leq 1\}$. Let $S \subseteq X$ with $|S| = s \leq s_0$. The number of vertices of Y adjacent to at least one vertex from S is the binomial random variable $Z \stackrel{d}{=} \text{Bi}(t, 1 - (1-p)^s)$, whose expected value is

$$(1 - (1-p)^s)t \geq (1 - e^{-ps})t \geq (1 - e^{-1})pst.$$

Hence, applying the Chernoff bound ([Lemma 2.1](#)), the set S fails the Hall condition with probability at most

$$\mathbb{P}(Z < s) = \mathbb{P}\left[Z < \frac{\mathbb{E}[Z]}{(1 - e^{-1})pt}\right] \leq \exp\left[-\mathbb{E}[Z] \cdot H\left(\frac{1}{(1 - e^{-1})pt}\right)\right] \leq e^{-stp/100},$$

where $H(x) = x \ln x - x + 1$ and we have used that $\lim_{x \downarrow 0} H(x) = 1$ and the last inequality holds for t sufficiently large. Hence, the probability that the necessary condition in the statement of Hall's theorem fails for at least one set S with $|S| \leq s_0$ is at most

$$\sum_{s=1}^{s_0} \binom{t}{s} e^{-stp/100} \leq \sum_{s=1}^{s_0} t^s e^{-stp/100} = \sum_{s=1}^{s_0} (te^{-tp/100})^s = O(te^{-tp/100}),$$

using that $te^{-tp/100} = o(1)$. Now, let $0 < \varepsilon < (1 - e^{-1})/2$ be a constant so that $\sum_{s \leq \varepsilon t} \binom{t}{s} \leq \exp(0.08t)$. Consider any set $S \subseteq X$ with $s_0 < |S| = s \leq \varepsilon t$. The expected size of $N[S]$ is at least $(1 - e^{-1})t$. It follows from Chernoff bound (see [Lemma 2.1](#)) that

$$\mathbb{P}\left(|N[S]| \leq (1 - e^{-1})t/2\right) \leq \exp[-H(1/2) \cdot (1 - e^{-1})t] \leq e^{-0.09t},$$

where $H(x) = x \ln x - x + 1$ is again as in [Lemma 2.1](#). The probability that the necessary condition fails for at least one set S with $s_0 < |S| \leq \varepsilon t$ is therefore at most

$$\sum_{s=s_0+1}^{\varepsilon t} \binom{t}{s} e^{-0.09t} \leq e^{-0.01t} = e^{-\Omega(tp)}.$$

Finally, let $S \subseteq X$ with $\varepsilon t < |S| = s \leq t$. If S fails the test, then there exists $T \subseteq Y$ of cardinality $t - s + 1$ such that there is no edge between S and T . Hence, the probability that the condition fails for at least one set S with $\varepsilon t < |S| \leq t$ is at most

$$\begin{aligned} \sum_{\varepsilon t < s \leq t} \binom{t}{s} \binom{t}{t-s+1} (1-p)^{s(t-s+1)} &\leq \sum_{\varepsilon t < s \leq t} t^{t-s} \cdot t^{t-s+1} \cdot \exp[-ps(t-s+1)] \\ &\leq \sum_{\varepsilon t < s \leq t} \exp[(t-s+1)2 \ln t - \varepsilon pt(t-s+1)] \\ &= \sum_{\varepsilon t < s \leq t} \exp[(t-s+1)(2 \ln t - \varepsilon pt)] \\ &\leq \sum_{\varepsilon t < s \leq t} \exp[-(t-s+1)\varepsilon pt/2] \\ &= O(e^{-\varepsilon pt/2}), \end{aligned}$$

where we have used that we can choose K large enough for $2 \ln t < \varepsilon K \ln t/2$ to hold in the penultimate line. We conclude that, using Hall's theorem:

$$\begin{aligned} \mathbb{P}(\mathcal{B}(t, p) \text{ has a perfect matching}) &= 1 - O(te^{-tp/100}) - e^{-0.09tp} - O(e^{-\varepsilon pt/2}) \\ &= 1 - O(te^{-\Omega(tp)}), \end{aligned}$$

as required. \square

Before we can proceed with our upper bound on the acquaintance time of the percolated random geometric graph we need some more preparations. A precise result on the acquaintance time of the binomial random graph $\mathcal{G}(n, p)$ was already given in [\[17\]](#), but we again require a version with precise quantitative bounds on the error probabilities. The following result will serve our purposes.

Theorem 6.2. *There exists a constant $\gamma > 0$ such that for all $k \leq t/1000$:*

$$\mathbb{P}\left[\mathcal{AC}(\mathcal{G}(t, p)) \leq k\right] \geq 1 - t^2 e^{-\gamma pk},$$

for all $t \in \mathbb{N}$ and $0 < p < 1$.

Proof. In order to avoid technical problems with events not being independent, we use a classic technique known as *multi-round exposure*. (In fact, we will use a three-round exposure here.) The observation is that a random graph $G \in \mathcal{G}(t, p)$ can be viewed as a union of three independently generated random graphs $G_1, G_2, G_3 \in \mathcal{G}(t, \bar{p})$, with \bar{p} defined by:

$$p = 1 - (1 - \bar{p})^3.$$

(See, for example, [\[6,14\]](#) for more information). Let us observe that $p \geq \bar{p} \geq p/3$.

Firstly, let us focus on $G_1 = (V, E) \in \mathcal{G}(t, \bar{p})$. Our goal is to show that, with probability at least $1 - O(te^{-\Omega(tp)})$, G_1 contains a path P of length $0.9t$. We consider the following process. Select any

vertex $v_1 \in V$ and expose all edges in G_1 from v_1 to other vertices of V . If at least one edge is found, select any neighbour v_2 of v_1 and expose all edges from v_2 to $V \setminus \{v_1, v_2\}$ with a hope that at least one edge is discovered and the process can be continued. The only reason for the process to terminate at a given round is that no edge is found. The probability that the process does *not* stop before discovering a path P of length $0.9t$ is equal to

$$\begin{aligned} \mathbb{P}(G_1 \text{ contains a path of length } \geq 0.9t) &\geq (1 - (1 - \bar{p})^{t-1}) \cdots (1 - (1 - \bar{p})^{0.1t+1}) \\ &\geq (1 - e^{-0.1\bar{p}t})^{0.9t} \\ &\geq 1 - 0.9te^{-0.1\bar{p}t}. \end{aligned}$$

Now let $k \leq t/1000$ be arbitrary. Conditioning on the event that G_1 has at least one path of length $0.9t$, let us fix a path $P \subseteq G_1$ of length $V(P) \geq 0.8t$ such that $V(P)$ is a multiple of k . We now consider $G_2 = (V, E) \in \mathcal{G}(t, \bar{p})$. It follows from Lemma 6.1 that, with probability at least $1 - O(te^{-\Omega(tp)})$, there is a matching $M \subseteq G_2$ between $V(P)$ and $V \setminus V(P)$ that saturates $V \setminus V(P)$. We call agents occupying $V(P)$ *active* and agents occupying $V \setminus V(P)$ *inactive*.

We split the path P into many paths, each on exactly k vertices. This partition also divides the active agents into $O(t/k)$ teams, each team consisting of k agents. Every team performs (independently and simultaneously) the strategy from Lemma 2.6. This certainly results in every pair of active vertices on the same team getting acquainted.

Next, we will consider the probability that an active agent x gets acquainted with an agent y that is either on a different team or inactive. It follows from Lemma 2.6 that agent x visits k distinct vertices. Since y either belongs to a different team or is inactive, the pair x, y occupy at least k distinct pairs of vertices during the process. Considering only those edges in $G_3 \in \mathcal{G}(t, p)$, the probability that the two agents never got acquainted is at most

$$\mathbb{P}(x, y \text{ do not get acquainted}) \leq (1 - \bar{p})^k \leq e^{-\bar{p}k}.$$

Since there are at most $\binom{t}{2}$ pairs of agents, the union bound shows:

$$\mathbb{P}(\text{all active vertices get acquainted to each other and all inactive vertices}) \geq 1 - O(t^2 e^{-\bar{p}k}).$$

Finally, to also acquaint the inactive vertices with each other, we simply use the matching M to place them on P , and repeat the “teams” strategy. This way they all pairs will indeed get acquainted. Summarizing, we have

$$\begin{aligned} \mathbb{P}(\text{all pairs get acquainted}) &\geq 1 - 0.9te^{-0.1\bar{p}t} - O(te^{-\Omega(tp)}) - O(t^2 e^{-k\bar{p}}) \\ &= 1 - O(t^2 e^{-\Omega(k\bar{p})}), \end{aligned}$$

which concludes the proof. \square

Now, we are ready to come back to the upper bound for the acquaintance time of percolated random geometric graphs. We will prove the following upper bound

Lemma 6.3. *Let $\varepsilon > 0$ be arbitrary. There is a constant $K > 0$ such that if the sequences $(r_n)_n$ and $(p_n)_n$ are such that $p_n < 1 - \varepsilon$ and $p_n n r_n^2 \geq K \ln n$ for all n then for $G \in \mathcal{G}(n, r_n, p_n)$ we have*

$$\mathcal{AC}(G) = O\left(\frac{\ln n}{r_n^2 \cdot p_n}\right) \quad \text{w.h.p.}$$

Proof. We want to mimic the strategy introduced for dense (classic) random geometric graph $\mathcal{G}(n, r_n)$ —see Section 4.1. Let us recall the setting briefly. \mathcal{D}_n denotes the dissection of the unit square into m_n^2 ($m_n := \lceil 1000/r_n \rceil$) equal squares called cells. For a given cell $c \in \Gamma_n$, $V(c)$ denotes the set of points of \mathcal{X}_n that fall in c ; $\mu_n := n/m_n^2$ is the expectation $\mathbb{E}|V(c)|$ and by Lemma 4.1, w.h.p., $0.9 \cdot \mu_n \leq |V(c)| \leq 1.1 \cdot \mu_n$ for every cell. Again each $c \in \Gamma_n$, $V(c)$ is partitioned into three parts $V_1(c), V_2(c), V_3(c)$, each of cardinality at most $0.4 \cdot \mu_n$. For each pair $1 \leq i < j \leq 3$ and each cell $c \in \Gamma_n$, $W_{ij}(c) \subseteq V(c)$ is a set of cardinality exactly

$$t := \lfloor 0.9\mu_n \rfloor,$$

such that $V_i(c) \cup V_j(c) \subseteq W_{ij}(c)$; and set $W_{ij} := \bigcup_{c \in \mathcal{I}_n} W_{ij}(c)$. $G_{ij} = G[W_{ij}]$ denotes the subgraph induced by W_{ij} .

Let E_{ij} denote the event that for each adjacent pair of cells $c, d \in \mathcal{D}_n$ there is a perfect matching between $W_{ij}(c)$ and $W_{ij}(d)$. It follows from [Lemma 6.1](#) that

$$\mathbb{P}(E_{ij}) \geq 1 - 4m_n^2 t e^{-\Omega(tp_n)} \geq 1 - 4n^4 e^{-\Omega(p_n n r_n^2)} = 1 - o(1),$$

where the last inequality holds since $p_n n r_n^2 \geq K \ln n$ with K a sufficiently large constant.

Now let us set

$$k := \min \left(\frac{C \cdot \ln n}{p_n}, \frac{t}{1000} \right)$$

with $C > 0$ a constant to be chosen more precisely later, and let F_{ij} denote the event that for every adjacent pair of cells $c, d \in \mathcal{D}_n$ we have that $\mathcal{AC}(G[W_{ij}(c) \cup W_{ij}(d)]) \leq k$. It follows from [Theorem 6.2](#) that

$$\mathbb{P}(F_{ij}) \geq 1 - m_n^2 \cdot O(t^2 e^{-\Omega(k p_n)}) = 1 - O(n^4 e^{-\Omega(\min(C \ln n, p_n n r_n^2))}) = 1 - o(1),$$

using that $p_n n r_n^2 \geq K \ln n$ and that we can assume $C, K > 0$ are sufficiently large for the last equality to hold.

We have seen that, w.h.p., E_{ij} and F_{ij} hold for each $1 \leq i < j \leq 3$. If E_{ij}, F_{ij} both hold, we have the following strategy to acquaint all agents on W_{ij} . Again we treat the set of agents initially on a cell $c \in \mathcal{D}_n$ as a group. We move these groups from cell to cell following the strategy of [Theorem 2.3](#), where for each move of the original strategy for the $m \times m$ grid, we use the perfect matchings (that exist because E_{ij} holds) between adjacent cells to transfer entire groups. After each move of the grid strategy, we do the following. Observe that the grid can be covered by four matchings M_1, \dots, M_4 (take for instance horizontal/vertical edges with odd/even x/y -coordinates of the leftmost/top point). For each such matching M_i we do the following simultaneously for each of its edges $cd \in M_i$. We acquaint the groups of agents currently on $W_{ij}(c) \cup W_{ij}(d)$ and return them to their starting points in $2k$ moves (this can be done since F_{ij} holds, and using [Lemma 2.5](#)). Repeating this procedure for each of W_{12}, W_{13}, W_{23} , it is clear that this way we do acquaint all agents with each other in at most

$$O(m_n^2 \cdot k) = O\left(\frac{\ln n}{r_n^2 \cdot p_n}\right),$$

moves, as required. \square

7. Conclusion and further work

In this article, we have determined the likely value of the acquaintance time of random geometric graphs up to the leading constant, whenever the graph is w.h.p. connected. A very natural question is thus to also find the leading constant (if it even exists).

For percolated random geometric graphs we have also found the likely value of the acquaintance time, but we needed a slightly stronger assumption on the sequences r_n, p_n . Namely, these parameters needed to be chosen in such a way that all degrees are slightly larger than \sqrt{n} . On the other hand we were able to provide an upper bound that works close to the connectivity threshold for percolated graphs. Another natural open problem is thus to determine the likely value of the acquaintance time of percolated random geometric graphs when $p_n n r_n^2 = \Omega(\ln n)$ and $p_n n r_n^2 = O(n^{1/2+o(1)})$.

And, of course it would again be nice to have more detailed asymptotics for the acquaintance time of percolated random geometric graphs.

Acknowledgements

The authors gratefully acknowledge support from NSERC and Ryerson University.

Appendix A. Proof (sketch) of Lemma 4.2

We briefly sketch how the proof of Lemma 3.1 in [5] should be adapted to yield a proof of our Lemma 4.2. The main adaptation that is needed is that their Lemma 3.1 needs to be altered to allow for values of T that are a small constant times $\ln n$, as follows:

Lemma A.1. *Let $\eta, K, \varepsilon > 0$ be arbitrary but fixed, and let m be given by (4) and let \mathcal{X}_n be as in (2). There exists a $\delta = \delta(\eta, K, \varepsilon) > 0$ such that for every $T \leq \delta \cdot \ln n$, w.h.p., the following hold:*

- (i) *Out of every $K \times K$ block of cells, the area of the bad cells inside the block is at most $(1 + \varepsilon) \ln n/n$;*
- (ii) *Out of every $K \times K$ block of cells touching the boundary of the unit square, the area of the bad cells inside the block is at most $(1 + \varepsilon) \ln n/2n$.*
- (iii) *Every $K \times K$ block of cells touching a corner contains only good cells.*

The proof of Lemma 3.1 in [5] relies only on the Chernoff bound and it is easily seen that an (almost) verbatim repeat of the proof will prove the above lemma. The rest of the proof of our Lemma 4.2 can now follow the proof of Lemma 3.3 from [5] (almost) verbatim since the proof given there does not rely on any probabilistic calculations (only geometric ones), besides their Lemma 3.1.

Appendix B. The proof of Lemma 4.3

Before we can start describing the proof of Lemma 4.3, we need to cover some more background and notation. The usual random geometric graph $G(n, r) = G(\mathcal{X}_n; r)$ is sometimes also called the *binomial random geometric graph*. It is often useful to switch to a “Poissonized” version of the random geometric graph. By this we mean the following. Let $N_n \stackrel{d}{=} \text{Po}(n)$ be independent of X_1, X_2, \dots , and set (following [22]):

$$\mathcal{P}_n := \{X_1, \dots, X_{N_n}\}. \quad (6)$$

(Thus \mathcal{P}_n is a Poisson process with intensity n on the unit square and intensity 0 elsewhere—see for instance [16] for background on Poisson processes.) The *Poisson random geometric graph* is defined as $G_{\mathcal{P}}(n, r) := G(\mathcal{P}_n; r)$. We will make use of the following result from [5], which is nearly identical to Theorem 1.6 of [22].

Theorem B.1 ([5]). *Let \mathcal{P}_n be as in (6), and let $h(a_1, \dots, a_k; A)$ be a bounded measurable function defined on all tuples $(a_1, \dots, a_k; A)$ with $A \subseteq \mathbb{R}^2$ finite and $a_1, \dots, a_k \in A$. Let us write*

$$Z := \sum_{\substack{a_1, \dots, a_k \in \mathcal{P}_n, \\ a_1, \dots, a_k \text{ distinct}}} h(a_1, \dots, a_k; \mathcal{P}_n).$$

Then

$$\mathbb{E}Z = n^k \cdot \mathbb{E}h(Y_1, \dots, Y_k; \{Y_1, \dots, Y_k\} \cup \mathcal{P}_n),$$

where Y_1, \dots, Y_k are i.i.d. uniform on the unit square, and are independent of \mathcal{P}_n .

We are now ready for the proof of Lemma 4.3, which follows that of Lemma 3.5 in [5].

Proof of Lemma 4.3. Let us first note that if we increase the parameter r then some obstructions might cease to be obstructions, while vertices that were safe for the smaller r will still be safe. So in particular, it suffices to prove the statement for $r_n := \sqrt{\ln n / \pi n}$.

Next, let us remark that, by Lemmas 4.2 and A.1, no point of any obstruction is within $100r$ of a corner of the unit square. We will use an appropriate first moment argument to bound the number of obstructions that are not close to a corner and have too few crucial vertices. For the moment, let us fix $a, c \leq K \ln n$ where $K := \lceil \pi / \eta^2 \rceil \cdot C$ with $C = C(\eta)$ as in Lemma 4.4. (So in particular, w.h.p., the maximum degree is at most $K \ln n$. And hence, w.h.p., every obstruction has at most $K \ln n$ points as well.) It is convenient to switch to the Poissonized version of the random graph for the moment, first show there are no such obstructions in the Poissonized setting, and then deduce the result for the original setting. Let us say that u, v is an (a, c) -pair (with respect to r, V) if:

- (pr-1) $\|u - v\| < r/100$;
 (pr-2) $B(u, r - \|u - v\|) \setminus B(u, \|u - v\|)$ contains exactly a points of $V \setminus \{u, v\}$;
 (pr-3) $B(u, \|u - v\|)$ contains exactly c points of $V \setminus \{u, v\}$.

To motivate this definition, note that if A is an obstruction, and we take $u, v \in A$ a pair realizing $\text{diam}(A)$, then (u, v) will be an (a, c) -pair with $c = s - 2$. If, furthermore (str-1)–(str-5) hold, then each of the a points in $B(u, r - \|u - v\|)$ must be crucial. (Each such point must be safe, otherwise it would have been part of the obstruction.)

Let $R_{\text{sde}}^{(a,c)}$ denote the number of (a, c) -pairs (X_i, X_j) in \mathcal{P}_n for which X_i is within $100r$ of the boundary of $[0, 1]^2$, but not within $100r$ of a corner; let $R_{\text{mdl}}^{(a,c)}$ denote the number of (a, c) -pairs (X_i, X_j) for which X_i is more than $100r$ away from the boundary of $[0, 1]^2$, and set $R^{(a,c)} = R_{\text{sde}}^{(a,c)} + R_{\text{mdl}}^{(a,c)}$.

For $0 < z < r/100$ let us write:

$$\mu_1(z) := n \cdot \text{area}(B(u; r - z) \setminus B(u; z)),$$

$$\mu_3(z) := n \cdot \text{area}(B(u; z)),$$

where $u, v \in \mathbb{R}^2$ are two points with $\|u - v\| = z$.

Let us now consider R_{mdl} . Using [Theorem B.1](#) we find

$$\begin{aligned} \mathbb{E}R_{\text{mdl}}^{(a,c)} &= n^2 \int_{[100r, 1-100r]^2} \int_{B(v; r/100)} \frac{\mu_1(\|u - v\|)^a e^{-\mu_1(\|u - v\|)}}{a!} \cdot \frac{\mu_3(\|u - v\|)^c e^{-\mu_3(\|u - v\|)}}{c!} \, du \, dv \\ &= n^2 (1 - 200r)^2 \int_0^{r/100} \frac{\mu_1(z)^a e^{-\mu_1(z)}}{a!} \cdot \frac{\mu_3(z)^c e^{-\mu_3(z)}}{c!} 2\pi z \, dz \\ &\leq n^2 \int_0^{r/100} (\pi n r^2)^a (\pi n z^2)^c e^{-\pi n r^2 - \pi n z^2} 2\pi z \, dz \\ &= O\left(n^{2+c} \cdot \ln^a n \int_0^{r/100} e^{-\ln n - \pi n z^2} z^{2c+1} \, dz\right) \\ &= O\left(n^{1+c} \cdot \ln^a n \int_0^{r/100} e^{-\pi n z^2} z^{2c+1} \, dz\right) \\ &= O\left(n^{1+c} \cdot \ln^a n \cdot (nr)^{-(2+2c)}\right) \\ &= O\left(\ln^a n \cdot (nr^2)^{-(1+c)}\right) \\ &= O\left((\ln n)^{a-(c+1)}\right). \end{aligned} \tag{7}$$

Here we have used a switch to polar coordinates to get the second line; and the change of variables $y = \pi n z^2$ to get the sixth line.

Now we turn attention to $R_{\text{sde}}^{(a,c)}$. By [Theorem B.1](#) and a switch to polar coordinates and where $z = \|u - v\|$ and w is the distance of the nearest of u, v to the boundary:

$$\begin{aligned} \mathbb{E}R_{\text{sde}}^{(a,b)} &\leq 800n^2 r \int_0^{r/100} (\pi n r^2)^a (\pi n z^2)^c e^{-\frac{\pi}{2} n r^2 - \frac{1}{2} \pi n z^2} \pi z \, dz \\ &= O\left(n^{2+c} \cdot r \cdot \ln^a n \int_0^{r/100} z^{2c+1} e^{-\frac{\pi}{2} n r^2 - \frac{1}{2} \pi n z^2} \, dz\right) \\ &= O\left(n^{\frac{3}{2}+c} \cdot \ln^a n \int_0^{r/100} z^{2c+1} e^{-\frac{1}{2} \pi n z^2} \, dz\right) \\ &= O\left(n^{\frac{3}{2}+c} \cdot r \cdot \ln^a n \cdot (nr)^{-(2c+2)}\right) \\ &= O\left(\ln^a n \cdot (nr^2)^{-(\frac{3}{2}+c)}\right) \\ &= O\left((\ln n)^{a-\frac{3}{2}-c}\right). \end{aligned}$$

We find that, for all $a \leq K \ln n$, $a + 100 \leq c \leq K \ln n$:

$$\mathbb{P}(R^{(a,c)} > 0) = O(\ln^{-100} n), \quad (8)$$

This is of course still in the Poissonized setting. Let $\tilde{R}^{(a,c)}$ be the corresponding number of pairs wrt. \mathcal{X}_n , r . (I.e. the number of (a, c) -pairs in \mathcal{X}_n that are not close to a corner.) Let us write $k := a + c + 2$ for notational convenience. If we condition on the event that $\tilde{R}^{(a,c)} > 0$ then we can sample a sequence $(Y_1, \dots, Y_k) \subseteq (\mathcal{X}_n)^k$, uniformly at random from all such sequences where (Y_1, Y_2) is an (a, c) -pair, $Y_3, \dots, Y_{c+2} \in B(Y_1, \|Y_1 - Y_2\|)$ and $Y_{c+3}, \dots, Y_k \in B(Y_1, r - \|Y_1 - Y_2\|)$. By symmetry considerations, every sequence without repetitions from $(\mathcal{X}_n)^k$ will be equally likely. Let E denote the event that $Y_1, \dots, Y_k \in \{X_i : i \leq n - n^{.99}\}$. The symmetry considerations show that:

$$\mathbb{P}(E^c | \tilde{R}^{(a,c)} > 0) \leq k \cdot n^{-.01} = o(1).$$

(We condition on $\tilde{R} > 0$ so that Y_1, Y_2, \dots, Y_k are defined.) Let F denote the event that $n - n^{.99} < N_n \leq n$. Since N_n is independent of the sequence of points X_1, X_2, \dots , the event F is independent of the events E and $\{R' > 0\}$. Hence we also have

$$\mathbb{P}(E^c | \tilde{R}^{(a,c)} > 0, F) = o(1). \quad (9)$$

Next, let us observe that if E and F hold then we must have that $R^{(a,c)} > 0$. (The points Y_1, Y_2, \dots, Y_k will all be part of \mathcal{P}_n as well as \mathcal{P}_n and since \mathcal{P}_n is contained in \mathcal{X}_n there are no additional points that could mess up the values of a or c by falling in the wrong region.) It follows that

$$\begin{aligned} \mathbb{P}(R^{(a,c)} > 0) &\geq \mathbb{P}(E | \tilde{R}^{(a,c)} > 0, F) \cdot \mathbb{P}(\tilde{R}^{(a,c)} > 0, F) \\ &= \mathbb{P}(E | \tilde{R}^{(a,c)} > 0, F) \cdot \mathbb{P}(\tilde{R}^{(a,c)} > 0) \cdot \mathbb{P}(F) \\ &= (1 - o(1)) \cdot \mathbb{P}(\tilde{R}^{(a,c)} > 0) \cdot \left(\frac{1}{2} - o(1)\right), \end{aligned} \quad (10)$$

using (9) and the Chernoff bound for the last line. Combining (8) and (10), it follows that we must have

$$\mathbb{P}(\tilde{R}^{(a,c)} > 0) = O(\ln^{-100} n),$$

as well. Setting $\tilde{R} := \sum_{\substack{a \leq K \ln n \\ a+100 \leq c \leq K \ln n}} \tilde{R}^{(a,c)}$, we see that we have $\tilde{R} = o(1)$ w.h.p. This finishes the proof that (w.h.p.) there cannot be any s -obstruction with less than $s - 100$ crucial vertices. \square

References

- [1] N. Alon, F.R.K. Chung, R.L. Graham, Routing permutations on graphs via matchings, *SIAM J. Discrete Math.* 7 (1994) 513–530.
- [2] O. Angel, I. Shinkar, A tight upper bound on acquaintance time of graphs, preprint. Available from <http://arxiv.org/abs/1307.6029>.
- [3] J. Balogh, B. Bollobás, M. Krivelevich, T. Müller, M. Walters, Hamilton cycles in random geometric graphs, *Ann. Appl. Probab.* 21 (3) (2011) 1053–1072.
- [4] I. Benjamini, I. Shinkar, G. Tsur, Acquaintance time of a graph, *SIAM J. Discrete Math.* 28 (2) (2014) 767–785.
- [5] A. Beveridge, A. Dudek, A.M. Frieze, T. Müller, M. Stojakovic, Maker-breaker games on random geometric graphs, *Random Structures Algorithms* 45 (4) (2014) 553–607.
- [6] B. Bollobás, *Random Graphs*, second ed., in: Cambridge Studies in Advanced Mathematics, vol. 73, Cambridge University Press, Cambridge, 2001.
- [7] N. Chen, On the approximability of influence in social networks, *SIAM J. Discrete Math.* 23 (5) (2009) 1400–1415.
- [8] C. Cooper, A. Frieze, The cover time of random geometric graphs, *Random Structures Algorithms* 38 (3) (2011) 324–349.
- [9] J. Czyzowicz, D. Dereniowski, L. Gasieniec, R. Klasing, A. Kosowski, D. Pajak, Collision-free network exploration, in: A. Pardo, A. Viola (Eds.), *LATIN 2014: Theoretical Informatics—11th Latin American Symposium*, Montevideo, Uruguay, March 31–April 4, 2014. Proceedings, in: *Lecture Notes in Computer Science*, vol. 8392, Springer, 2014, pp. 342–354.
- [10] A. Dudek, P. Prałat, Acquaintance time of random graphs near connectivity threshold, preprint.
- [11] E.N. Gilbert, Random plane networks, *J. Soc. Ind. Appl. Math.* 9 (1961) 533–543.
- [12] P. Gupta, P.R. Kumar, Critical power for asymptotic connectivity in wireless networks, in: *Stochastic Analysis, Control, Optimization and Applications*, in: *Systems Control Found. Appl.*, Birkhäuser, Boston, MA, 1999, pp. 547–566.
- [13] S.T. Hedetniemi, S.M. Hedetniemi, A. Liestman, A survey of gossiping and broadcasting in communication networks, *Networks* 18 (4) (1998) 319–349.

- [14] S. Janson, T. Łuczak, A. Ruciński, Random graphs, in: Wiley–Interscience Series in Discrete Mathematics and Optimization, Wiley–Interscience, New York, 2000.
- [15] D. Kempe, J. Kleinberg, E. Tardos, Maximizing the spread of influence through a social network, in: KKD, 2003, pp. 137–146.
- [16] J. Kingman, Poisson Processes, Oxford University Press, Oxford, 1993.
- [17] W. Kinnersley, D. Mitsche, P. Prałat, A note on the acquaintance time of random graphs, *Electron. J. Combin.* 20 (3) (2013) #P52.
- [18] G. Mao, B. D. O Anderson, Towards a better understanding of large-scale network models, *IEEE/ACM Trans. Netw.* 20 (2012) 408–421.
- [19] C.J.H. McDiarmid, T. Müller, On the chromatic number of random geometric graphs, *Combinatorica* 31 (4) (2011) 423–488.
- [20] T. Müller, X. Pérez-Giménez, N. Wormald, Disjoint Hamilton cycles in the random geometric graph, *J. Graph Theory* 68 (4) (2011) 299–322.
- [21] T. Müller, Two-point concentration in random geometric graphs, *Combinatorica* 28 (5) (2008) 529–545.
- [22] M.D. Penrose, *Random Geometric Graphs*, Oxford University Press, Oxford, 2003.
- [23] M.D. Penrose, The longest edge of the random minimal spanning tree, *Ann. Appl. Probab.* 7 (2) (1997) 340–361.
- [24] M.D. Penrose, On k -connectivity for a geometric random graph, *Random Structures Algorithms* 15 (2) (1999) 145–164.
- [25] M.D. Penrose, Connectivity of soft random geometric graphs, preprint, available from <http://arxiv.org/abs/1311.3897>.
- [26] D. Reichman, New bounds for contagious sets, *Discrete Math.* 312 (2012) 1812–1814.
- [27] O. Yağan, Performance of the Eschenauer–Gligor key distribution scheme under an ON/OFF channel, *IEEE Trans. Inform. Theory* 58 (6) (2012) 3821–3835.